

CSE 410/510 Special Topics: Software Security

Instructor: Dr. Ziming Zhao

Location: Obrian 109

Time: Monday, Wednesday 5:00PM-6:20PM

Last Class

1. Stack-based buffer overflow
 - a. Overwrite Saved EBP.

This Class

1. Stack-based buffer overflow
 - a. Defense.

crackme4h

```
void printsecret(int i, int j, int k)
{
    if (i == 0xdeadbeef && j == 0xCODECAFE && k == 0xD0D0FACE)
        print_flag();

    exit(0);}

int main(int argc, char *argv[])
{
    char buf[8];

    if (argc != 2)
        return 0;

    strcpy(buf, argv[1]);
}
```

crackme4

```
000012b7 <main>:
 12b7:f3 0f 1e fb      endbr32
12bb:55      push %ebp
12bc:89 e5      mov  %esp,%ebp
12be:83 ec 08    sub  $0x8,%esp
12c1:83 7d 08 02    cmpl $0x2,0x8(%ebp)
12c5:74 07      je   12ce <main+0x17>
12c7:b8 00 00 00 00    mov  $0x0,%eax
12cc:eb 1a      jmp  12e8 <main+0x31>
12ce:8b 45 0c    mov  0xc(%ebp),%eax
12d1:83 c0 04    add  $0x4,%eax
12d4:8b 00     mov  (%eax),%eax
12d6:50      push %eax
12d7:8d 45 f8    lea -0x8(%ebp),%eax
12da:50      push %eax
12db:e8 fc ff ff    call 12dc <main+0x25>
12e0:83 c4 08    add  $0x8,%esp
12e3:b8 00 00 00 00    mov  $0x0,%eax
12e8:c9      leave
12e9:c3      ret
12ea:66 90     xchg %ax,%ax
12ec:66 90     xchg %ax,%ax
12ee:66 90     xchg %ax,%ax
```

Arg3 = 0xd0doface

Arg2 = 0xcodecafe

Arg1 = 0xdeadbeef

4 bytes

RET = printsecret

crackme4h

000012c6 <main>:

12c6: f3 0f 1e fb endbr32

12ca: 8d 4c 24 04 lea 0x4(%esp),%ecx

12ce: 83 e4 f0 and \$0xffffffff0,%esp

12d1: ff 71 fc pushl -0x4(%ecx)

12d4: 55 push %ebp

12d5: 89 e5 mov %esp,%ebp

12d7: 51 push %ecx

12d8: 83 ec 14 sub \$0x14,%esp

12db: 89 c8 mov %ecx,%eax

12dd: 83 38 02 cmpl \$0x2,(%eax)

12e0: 74 07 je 12e9 <main+0x23>

12e2: b8 00 00 00 00 mov \$0x0,%eax

12e7: eb 1d jmp 1306 <main+0x40>

12e9: 8b 40 04 mov 0x4(%eax),%eax

12ec: 83 c0 04 add \$0x4,%eax

12ef: 8b 00 mov (%eax),%eax

12f1: 83 ec 08 sub \$0x8,%esp

12f4: 50 push %eax

12f5: 8d 45 f0 lea -0x10(%ebp),%eax

12f8: 50 push %eax

12f9: e8 fc ff ff call 12fa <main+0x34>

12fe: 83 c4 10 add \$0x10,%esp

1301: b8 00 00 00 00 mov \$0x0,%eax

1306: 8b 4d fc mov -0x4(%ebp),%ecx

1309: c9 leave

130a: 8d 61 fc lea -0x4(%ecx),%esp

130d: c3 ret

crackme4h

```
000012c6 <main>:
 12c6: f3 0f 1e fb     endbr32
 12ca: 8d 4c 24 04     lea 0x4(%esp),%ecx
 12ce: 83 e4 f0       and $0xffffffff0,%esp
 12d1: ff 71 fc       pushl -0x4(%ecx)
 12d4: 55            push %ebp
 12d5: 89 e5         mov %esp,%ebp
 12d7: 51            push %ecx
 12d8: 83 ec 14      sub $0x14,%esp
 12db: 89 c8         mov %ecx,%eax
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 12e9: 8b 40 04      mov 0x4(%eax),%eax
 12ec: 83 c0 04      add $0x4,%eax
 12ef: 8b 00         mov (%eax),%eax
 12f1: 83 ec 08      sub $0x8,%esp
 12f4: 50            push %eax
 12f5: 8d 45 f0      lea -0x10(%ebp),%eax
 12f8: 50            push %eax
 12f9: e8 fc ff ff   call 12fa <main+0x34>
 12fe: 83 c4 10      add $0x10,%esp
1301: b8 00 00 00 00 mov $0x0,%eax
1306: 8b 4d fc      mov -0x4(%ebp),%ecx
1309: c9            leave
130a: 8d 61 fc      lea -0x4(%ecx),%esp
130d: c3            ret
```

%ecx →

%esp →

argv[1]

argv[0]

argc

RET

crackme4h

```
000012c6 <main>:
 12c6: f3 0f 1e fb      endbr32
 12ca: 8d 4c 24 04      lea -0x4(%esp),%ecx
 12ce: 83 e4 f0      and $0xffffffff0,%esp
 12d1: ff 71 fc      pushl -0x4(%ecx)
 12d4: 55          push %ebp
 12d5: 89 e5      mov %esp,%ebp
 12d7: 51          push %ecx
 12d8: 83 ec 14    sub $0x14,%esp
 12db: 89 c8      mov %ecx,%eax
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 12e7: eb 1d      jmp 1306 <main+0x40>
 12e9: 8b 40 04    mov 0x4(%eax),%eax
 12ec: 83 c0 04    add $0x4,%eax
 12ef: 8b 00      mov (%eax),%eax
 12f1: 83 ec 08    sub $0x8,%esp
 12f4: 50          push %eax
 12f5: 8d 45 f0    lea -0x10(%ebp),%eax
 12f8: 50          push %eax
 12f9: e8 fc ff ff    call 12fa <main+0x34>
 12fe: 83 c4 10    add $0x10,%esp
 1301: b8 00 00 00 00  mov $0x0,%eax
 1306: 8b 4d fc    mov -0x4(%ebp),%ecx
 1309: c9          leave
 130a: 8d 61 fc    lea -0x4(%ecx),%esp
 130d: c3          ret
```

%ecx →

%esp →

argv[1]

argv[0]

argc

RET

Size <= 16 bytes

crackme4h

```
000012c6 <main>:
 12c6: f3 0f 1e fb      endbr32
 12ca: 8d 4c 24 04      lea 0x4(%esp),%ecx
 12ce: 83 e4 f0        and $0xffffffff0,%esp
 12d1: ff 71 fc        pushl -0x4(%ecx)
 12d4: 55              push %ebp
 12d5: 89 e5           mov %esp,%ebp
 12d7: 51             push %ecx
 12d8: 83 ec 14       sub $0x14,%esp
 12db: 89 c8          mov %ecx,%eax
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 12f1: 83 ec 08       sub $0x8,%esp
 12f4: 50            push %eax
 12f5: 8d 45 f0       lea -0x10(%ebp),%eax
 12f8: 50            push %eax
 12f9: e8 fc ff ff    call 12fa <main+0x34>
 12fe: 83 c4 10       add $0x10,%esp
1301: b8 00 00 00 00  mov $0x0,%eax
1306: 8b 4d fc       mov -0x4(%ebp),%ecx
1309: c9            leave
130a: 8d 61 fc       lea -0x4(%ecx),%esp
130d: c3            ret
```

%ecx →

%esp →

argv[1]

argv[0]

argc

RET

Size <= 16 bytes

RET

crackme4h

```
000012c6 <main>:
12c6: f3 0f 1e fb      endbr32
12ca: 8d 4c 24 04      lea 0x4(%esp),%ecx
12ce: 83 e4 f0        and $0xffffffff0,%esp
12d1: ff 71 fc        pushl -0x4(%ecx)
12d4: 55             push %ebp
12d5: 89 e5         mov %esp,%ebp
12d7: 51             push %ecx
12d8: 83 ec 14       sub $0x14,%esp
12db: 89 c8         mov %ecx,%eax
12dd: 83 38 02       cmpl $0x2,(%eax)
12e0: 74 07         je 12e9 <main+0x23>
12e2: b8 00 00 00 00  mov $0x0,%eax
12e7: eb 1d         jmp 1306 <main+0x40>
12e9: 8b 40 04       mov 0x4(%eax),%eax
12ec: 83 c0 04       add $0x4,%eax
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12f5: 8d 45 f0       lea -0x10(%ebp),%eax
12f8: 50             push %eax
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12fe: 83 c4 10       add $0x10,%esp
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1309: c9             leave
130a: 8d 61 fc       lea -0x4(%ecx),%esp
130d: c3             ret
```

%ecx →

argv[1]

argv[0]

argc

RET

Size <= 16 bytes

RET

%ebp, %esp →

Saved EBP

crackme4h

```
000012c6 <main>:
 12c6: f3 0f 1e fb      endbr32
 12ca: 8d 4c 24 04      lea 0x4(%esp),%ecx
 12ce: 83 e4 f0         and $0xffffffff0,%esp
 12d1: ff 71 fc         pushl -0x4(%ecx)
 12d4: 55              push %ebp
 12d5: 89 e5           mov %esp,%ebp
 12d7: 51             push %ecx
 12d8: 83 ec 14       sub $0x14,%esp
 12db: 89 c8         mov %ecx,%eax
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 12e0: 74 07        je 12e9 <main+0x23>
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 12e9: 8b 40 04      mov 0x4(%eax),%eax
 12ec: 83 c0 04      add $0x4,%eax
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 1301: b8 00 00 00 00  mov $0x0,%eax
 1306: 8b 4d fc     mov -0x4(%ebp),%ecx
 1309: c9          leave
 130a: 8d 61 fc     lea -0x4(%ecx),%esp
 130d: c3          ret
```

%ecx →

%ebp →

%esp →

argv[1]

argv[0]

argc

RET

Size <= 16 bytes

RET

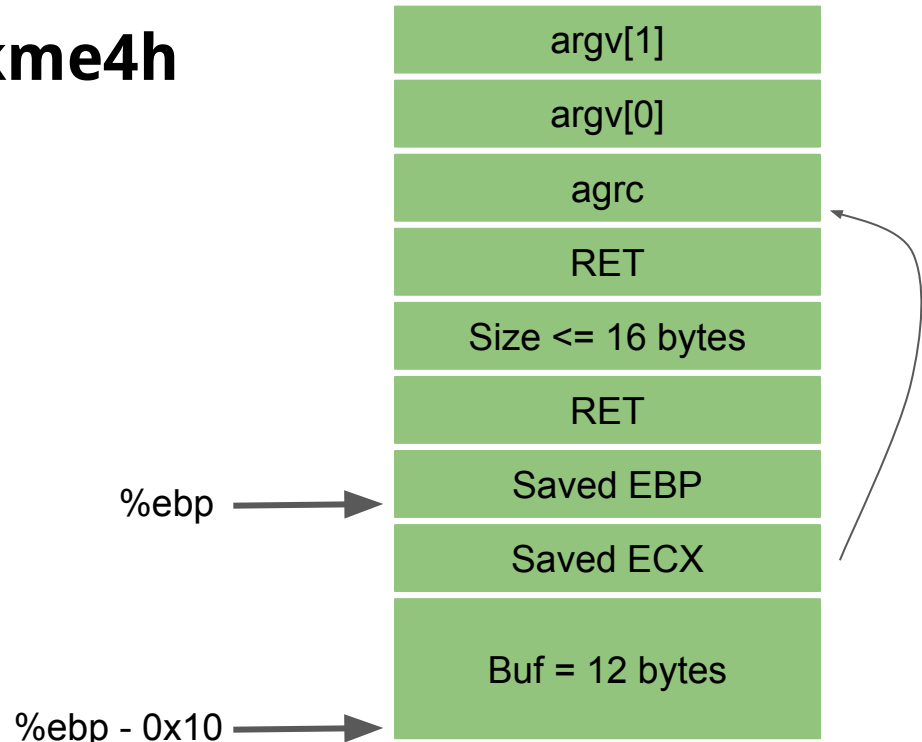
Saved EBP

Saved ECX



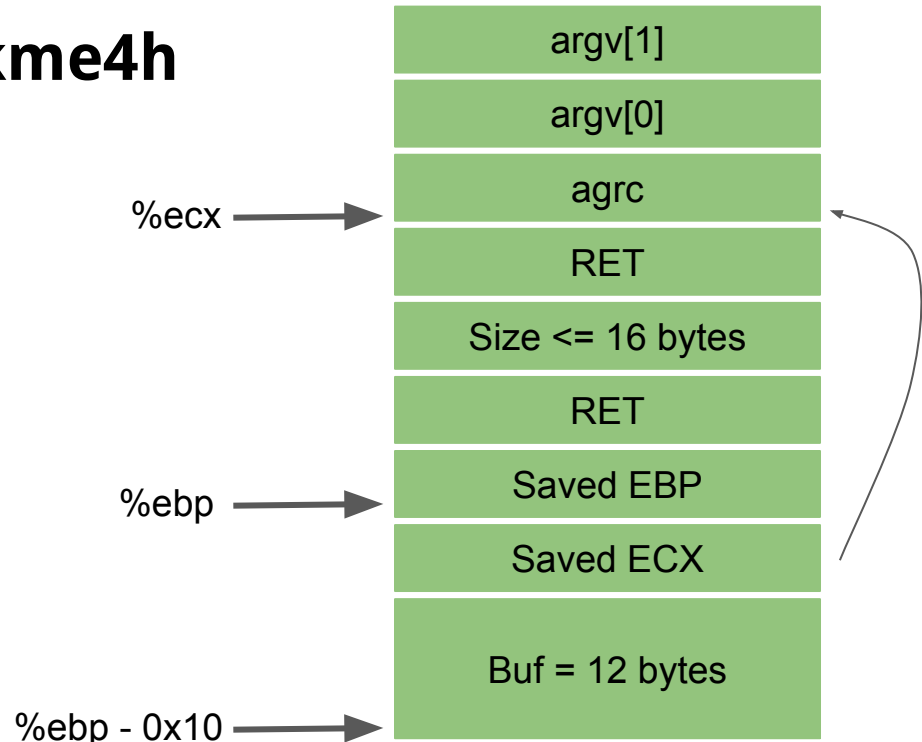
crackme4h

```
000012c6 <main>:
12c6: f3 0f 1e fb      endbr32
12ca: 8d 4c 24 04      lea 0x4(%esp),%ecx
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12d1: ff 71 fc         pushl -0x4(%ecx)
12d4: 55              push %ebp
12d5: 89 e5           mov %esp,%ebp
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12d8: 83 ec 14       sub $0x14,%esp
12db: 89 c8         mov %ecx,%eax
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12f4: 50          push %eax
12f5: 8d 45 f0      lea -0x10(%ebp),%eax
12f8: 50          push %eax
12f9: e8 fc ff ff ff call 12fa <main+0x34>
12fe: 83 c4 10      add $0x10,%esp
1301: b8 00 00 00 00  mov $0x0,%eax
1306: 8b 4d fc      mov -0x4(%ebp),%ecx
1309: c9          leave
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crackme4h

```
000012c6 <main>:
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130a: 8d 61 fc     lea -0x4(%ecx),%esp
130d: c3          ret
```

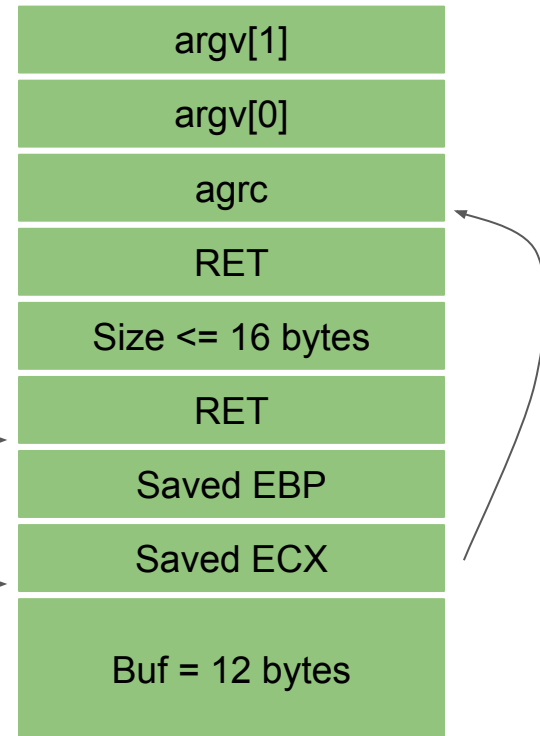


crackme4h

```
000012c6 <main>:
 12c6: f3 0f 1e fb      endbr32
 12ca: 8d 4c 24 04      lea 0x4(%esp),%ecx
 12ce: 83 e4 f0        and $0xffffffff0,%esp
 12d1: ff 71 fc        pushl -0x4(%ecx)
 12d4: 55             push %ebp
 12d5: 89 e5          mov %esp,%ebp
 12d7: 51            push %ecx
 12d8: 83 ec 14      sub $0x14,%esp
 12db: 89 c8        mov %ecx,%eax
 12dd: 83 38 02      cmpl $0x2,(%eax)
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 12f8: 50          push %eax
 12f9: e8 fc ff ff    call 12fa <main+0x34>
 12fe: 83 c4 10      add $0x10,%esp
 1301: b8 00 00 00 00  mov $0x0,%eax
 1306: 8b 4d fc      mov -0x4(%ebp),%ecx
 1309: c9          leave
 130a: 8d 61 fc      lea -0x4(%ecx),%esp
 130d: c3          ret
```

%esp →

%ecx →

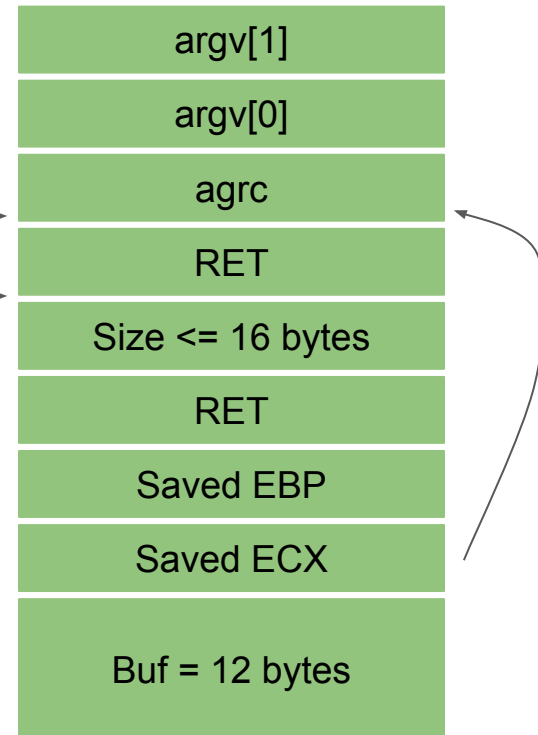


crackme4h

```
000012c6 <main>:
12c6: f3 0f 1e fb    endbr32
12ca: 8d 4c 24 04    lea 0x4(%esp),%ecx
12ce: 83 e4 f0      and $0xffffffff0,%esp
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12f4: 50       push %eax
12f5: 8d 45 f0    lea -0x10(%ebp),%eax
12f8: 50       push %eax
12f9: e8 fc ff ff  call 12fa <main+0x34>
12fe: 83 c4 10    add $0x10,%esp
1301: b8 00 00 00 00  mov $0x0,%eax
1306: 8b 4d fc    mov -0x4(%ebp),%ecx
1309: c9       leave
130a: 8d 61 fc    lea -0x4(%ecx),%esp
130d: c3       ret
```

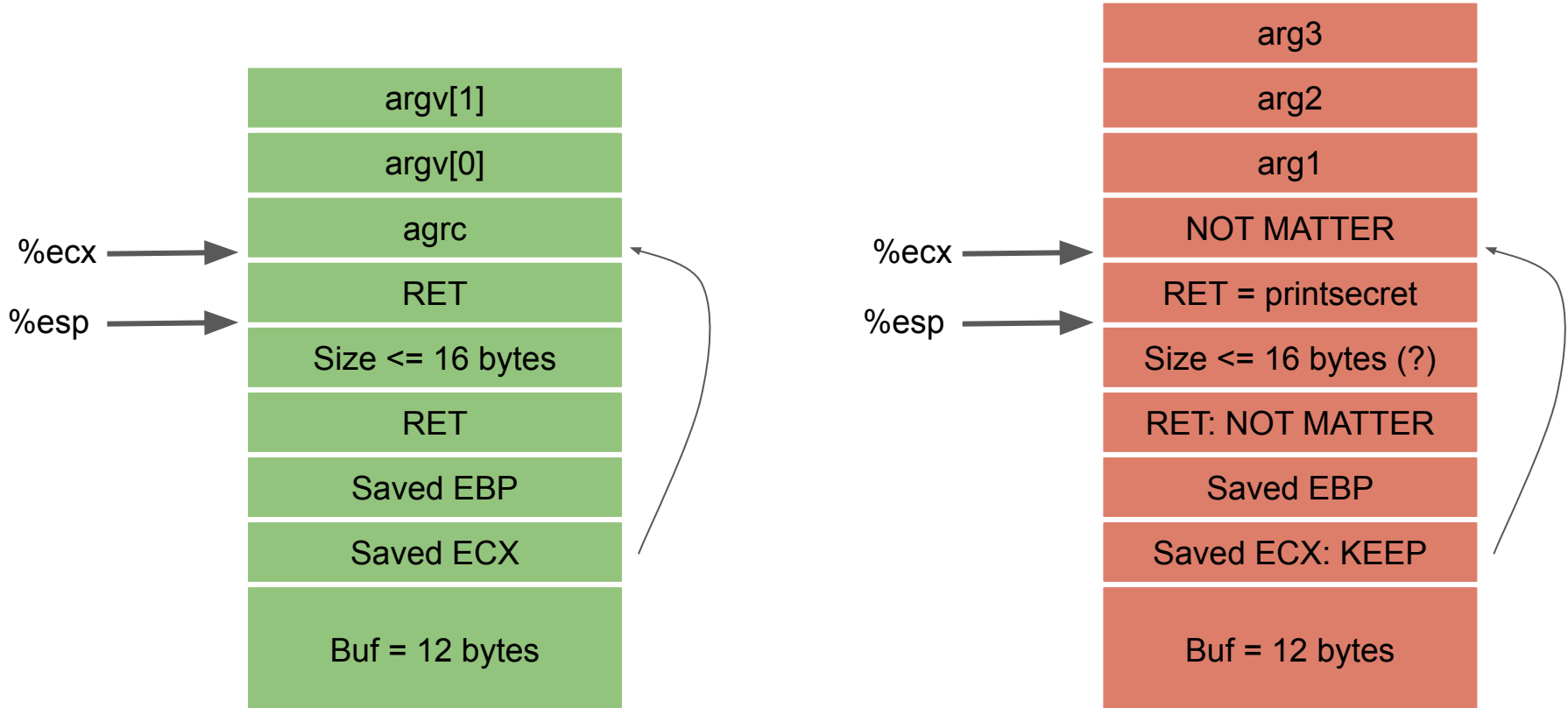
%ecx →

%esp →



Crackme4h

Craft the exploit



crackme464

```
0000000000001200 <printsecret>:
 1200:f3 0f 1e fa      endbr64
 1204:55                push %rbp
 1205:48 89 e5          mov  %rsp,%rbp
 1208:48 83 ec 10        sub  $0x10,%rsp
 120c:89 7d fc            mov  %edi,-0x4(%rbp)
 120f:89 75 f8          mov  %esi,-0x8(%rbp)
 1212:89 55 f4            mov  %edx,-0xc(%rbp)
 1215:81 7d fc ef be ad de  cmpl $0xdeadbeef,-0x4(%rbp)
 121c:75 32              jne  1250 <printsecret+0x50>
 121e:81 7d f8 fe ca de c0  cmpl $0xc0decafe,-0x8(%rbp)
 1225:75 29              jne  1250 <printsecret+0x50>
 1227:81 7d f4 ce fa d0 d0  cmpl $0xd0d0face,-0xc(%rbp)
 122e:75 20              jne  1250 <printsecret+0x50>
 1230:48 8d 3d d9 2d 00 00  lea  0x2dd9(%rip),%rdi   # 4010 <s>
 1237:e8 6d ff ff ff      callq 11a9 <decrypt>
 123c:48 89 c6            mov  %rax,%rsi
 123f:48 8d 3d c2 0d 00 00  lea  0xdc2(%rip),%rdi   # 2008 <_IO_stdin_used+0x8>
 1246:b8 00 00 00 00      mov  $0x0,%eax
 124b:e8 50 fe ff ff      callq 10a0 <printf@plt>
 1250:bf 00 00 00 00      mov  $0x0,%edi
 1255:e8 56 fe ff ff      callq 10b0 <exit@plt>
```

Return to here!!

Conditions we depend on to pull off the attack of *returning to shellcode on stack*

1. The ability to put the shellcode onto stack (env, command line)
2. The stack is executable
3. The ability to overwrite RET addr on stack before instruction **ret** is executed or to overwrite Saved EBP
4. Know the address of the destination function

Conditions we depend on to pull off the attack of *returning to shellcode on stack*

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- ~~2. The stack is executable~~
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**Defense 1:
Data Execution Prevention
(DEP, W \oplus X, NX)**

Harvard vs. Von-Neumann Architecture

Harvard Architecture

The Harvard architecture stores machine instructions and data in separate memory units that are connected by different busses. In this case, there are at least two memory address spaces to work with, so there is a memory register for machine instructions and another memory register for data. Computers designed with the Harvard architecture are able to run a program and access data independently, and therefore simultaneously. Harvard architecture has a strict separation between data and code. Thus, Harvard architecture is more complicated but separate pipelines remove the bottleneck that Von Neumann creates.

Von-Neumann architecture

In a Von-Neumann architecture, the same memory and bus are used to store both data and instructions that run the program. Since you cannot access program memory and data memory simultaneously, the Von Neumann architecture is susceptible to bottlenecks and system performance is affected.

Older CPUs

Older CPUs: Read permission on a page implies execution. So all readable memory was executable.

AMD64 – introduced NX bit (No-eXecute) in 2003

Windows Supporting DEP from Windows XP SP2 (in 2004)

Linux Supporting NX since 2.6.8 (in 2004)

Modern CPUs

Modern architectures support memory permissions:

- **PROT_READ** allows the process to read memory
- **PROT_WRITE** allows the process to write memory
- **PROT_EXEC** allows the process to execute memory

gcc parameter **-z *execstack*** to disable this protection


```
zining@zining-XPS-13-9300:~/Dropbox/myTeaching/System Security - Attack and Defense for Binaries UB 2020/code/overflow6$ readelf -l of6
```

```
Elf file type is DYN (Shared object file)
```

```
Entry point 0x1090
```

```
There are 12 program headers, starting at offset 52
```

```
Program Headers:
```

| Type | Offset | VirtAddr | PhysAddr | FileSiz | MemSiz | Flg | Align |
|--|----------|------------|------------|---------|---------|-----|--------|
| PHDR | 0x000034 | 0x00000034 | 0x00000034 | 0x00180 | 0x00180 | R | 0x4 |
| INTERP | 0x0001b4 | 0x000001b4 | 0x000001b4 | 0x00013 | 0x00013 | R | 0x1 |
| [Requesting program interpreter: /lib/ld-linux.so.2] | | | | | | | |
| LOAD | 0x000000 | 0x00000000 | 0x00000000 | 0x003f8 | 0x003f8 | R | 0x1000 |
| LOAD | 0x001000 | 0x00001000 | 0x00001000 | 0x002d4 | 0x002d4 | R E | 0x1000 |
| LOAD | 0x002000 | 0x00002000 | 0x00002000 | 0x001ac | 0x001ac | R | 0x1000 |
| LOAD | 0x002ed8 | 0x00003ed8 | 0x00003ed8 | 0x00130 | 0x00134 | RW | 0x1000 |
| DYNAMIC | 0x002ee0 | 0x00003ee0 | 0x00003ee0 | 0x000f8 | 0x000f8 | RW | 0x4 |
| NOTE | 0x0001c8 | 0x000001c8 | 0x000001c8 | 0x00060 | 0x00060 | R | 0x4 |
| GNU_PROPERTY | 0x0001ec | 0x000001ec | 0x000001ec | 0x0001c | 0x0001c | R | 0x4 |
| GNU_EH_FRAME | 0x002008 | 0x00002008 | 0x00002008 | 0x0005c | 0x0005c | R | 0x4 |
| GNU_STACK | 0x000000 | 0x00000000 | 0x00000000 | 0x00000 | 0x00000 | RWE | 0x10 |
| GNU_RELRO | 0x002ed8 | 0x00003ed8 | 0x00003ed8 | 0x00128 | 0x00128 | R | 0x1 |

```
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```

```
Elf file type is DYN (Shared object file)
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```
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There are 12 program headers, starting at offset 52
```

```
Program Headers:
```

| Type | Offset | VirtAddr | PhysAddr | FileSiz | MemSiz | Flg | Align |
|--|----------|------------|------------|---------|---------|-----|--------|
| PHDR | 0x000034 | 0x00000034 | 0x00000034 | 0x00180 | 0x00180 | R | 0x4 |
| INTERP | 0x0001b4 | 0x000001b4 | 0x000001b4 | 0x00013 | 0x00013 | R | 0x1 |
| [Requesting program interpreter: /lib/ld-linux.so.2] | | | | | | | |
| LOAD | 0x000000 | 0x00000000 | 0x00000000 | 0x003f8 | 0x003f8 | R | 0x1000 |
| LOAD | 0x001000 | 0x00001000 | 0x00001000 | 0x002d4 | 0x002d4 | R E | 0x1000 |
| LOAD | 0x002000 | 0x00002000 | 0x00002000 | 0x001ac | 0x001ac | R | 0x1000 |
| LOAD | 0x002ed8 | 0x00003ed8 | 0x00003ed8 | 0x00130 | 0x00134 | RW | 0x1000 |
| DYNAMIC | 0x002ee0 | 0x00003ee0 | 0x00003ee0 | 0x000f8 | 0x000f8 | RW | 0x4 |
| NOTE | 0x0001c8 | 0x000001c8 | 0x000001c8 | 0x00060 | 0x00060 | R | 0x4 |
| GNU_PROPERTY | 0x0001ec | 0x000001ec | 0x000001ec | 0x0001c | 0x0001c | R | 0x4 |
| GNU_EH_FRAME | 0x002008 | 0x00002008 | 0x00002008 | 0x0005c | 0x0005c | R | 0x4 |
| GNU_STACK | 0x000000 | 0x00000000 | 0x00000000 | 0x00000 | 0x00000 | RW | 0x10 |
| GNU_RELRO | 0x002ed8 | 0x00003ed8 | 0x00003ed8 | 0x00128 | 0x00128 | R | 0x1 |

What DEP cannot prevent

Can still corrupt stack or function pointers or critical data on the heap

As long as RET (saved EIP) points into legit code section, W \oplus X protection will not block control transfer

Ret2libc 32bit Bypassing NX

Discovered by Solar Designer, 1997

Ret2libc

Now programs built with non-executable stack.

Then, how to run a shell? Ret to C library ***system("/bin/sh")*** like how we called `printsecret()` in `overflowret`

Description

The C library function `int system(const char *command)` passes the command name or program name specified by `command` to the host environment to be executed by the command processor and returns after the command has been completed.

Declaration

Following is the declaration for `system()` function.

```
int system(const char *command)
```

Parameters

- `command` – This is the C string containing the name of the requested variable.

Return Value

The value returned is `-1` on error, and the return status of the command otherwise.

Buffer Overflow Example: code/overflowret4 32-bit (overflowret4_no_excNry_32)

```
int vulfoo()
{
    char buf[30];

    gets(buf);
    return 0;
}

int main(int argc, char *argv[])
{
    vulfoo();
    printf("I pity the fool!\n");
}
```

Conditions we depend on to pull off the attack of *ret2libc*

- ~~1. The ability to put the shellcode onto stack (env, command line)~~
- ~~2. The stack is executable~~
3. The ability to overwrite RET addr on stack before instruction **ret** is executed or to overwrite Saved EBP
4. Know the address of the destination function and arguments

Control Hijacking Attacks

Control flow

- Order in which individual statements, instructions or function calls of a program are executed or evaluated

Control Hijacking Attacks (Runtime exploit)

- A control hijacking attack exploits a program error, particularly a memory corruption vulnerability, at application runtime to subvert the intended control-flow of a program.
- Alter a code pointer (i.e., value that influences program counter) or, Gain control of the instruction pointer `%eip`
- Change memory region that should not be accessed

Code Injection Attacks

Code-injection Attacks

- a subclass of control hijacking attacks that subverts the intended control-flow of a program to previously injected malicious code

Shellcode

- code supplied by attacker – often saved in buffer being overflowed – traditionally transferred control to a shell (user command-line interpreter)
- machine code – specific to processor and OS – traditionally needed good assembly language skills to create – more recently have automated sites/tools

Code-Reuse Attack

Code-Reuse Attack: a subclass of control-flow attacks that subverts the intended control-flow of a program to invoke an unintended execution path inside the original program code.

Return-to-Libc Attacks (Ret2Libc)

Return-Oriented Programming (ROP)

Jump-Oriented Programming (JOP)

Exercise: Overthewire /maze/maze2

Overthewire

<http://overthewire.org/wargames/>

1. Open a terminal
2. Type: `ssh -p 2225 maze2@maze.labs.overthewire.org`
3. Input password: `fooghihahr`
4. `cd /maze`; this is where the binary are
5. Your goal is to get the password of maze3