CSE 703 Seminar: Advanced Software Security - Techniques and Tools

Instructor: Dr. Ziming Zhao

Location: Online Time: Monday, 12:50 PM-2:55 PM

Writing LLVM Passes

Passes

- LLVM applies a chain of analyses and transformations to the target program
- Each of these analyses or transformations is called a **pass**
- Machine-independent passes are invoked by *opt*
- Machine-dependant passes are invoked by *llc*
- A pass may require information provided by other passes. Dependencies must be explicitly stated. A transformation pass may require an analysis pass.

Types of Passes

A pass is an instance of the LLVM class Pass.

LLVM Coding Basics

- Written in modern C++, uses the STL:
 - Particularly the vector, set, and map classes
- LLVM IR is almost all doubly-linked lists:
 - Module contains lists of Functions & GlobalVariables
 - Function contains lists of BasicBlocks & Arguments
 - BasicBlock contains list of Instructions
- Linked lists are traversed with iterators:
 Function *M = ...
 for (Function::iterator I = M->begin(); I != M->end(); ++I) {
 BasicBlock &BB = *I;

LLVM Pass Manager

- Compiler is organized as a series of 'passes':
 - Each pass is one analysis or transformation
- Four types of Pass:
 - ModulePass: general interprocedural pass
 - CallGraphSCCPass: bottom-up on the call graph
 - FunctionPass: process a function at a time
 - BasicBlockPass: process a basic block at a time
- Constraints imposed (e.g. FunctionPass):
 - FunctionPass can only look at "current function"
 - Cannot maintain state across functions

Services provided by PassManager

- Optimization of pass execution:
 - Process a function at a time instead of a pass at a time
 - Example: If F, G, H are three functions in input pgm: "FFFFGGGGGHHHH" not "FGHFGHFGHFGH"
 - Process functions in parallel on an SMP (future work)
- Declarative dependency management:
 - Automatically fulfill and manage analysis pass lifetimes
 - Share analyses between passes when safe:
 - e.g. "DominatorSet live unless pass modifies CFG"
- Avoid boilerplate for traversal of program

See also: docs/WritingAnLLVMPass.html

HelloWorld Pass

#include "llvm/Pass.h"
#include "llvm/IR/Function.h"
#include "llvm/Support/raw_ostream.h"

#include "llvm/IR/LegacyPassManager.h"
#include "llvm/Transforms/IPO/PassManagerBuilder.h"

using namespace llvm;

namespace {
struct Hello : public FunctionPass {
 static char ID;
 Hello() : FunctionPass(ID) {}

bool runOnFunction(Function &F) override {
 errs() << "Hello: ";
 errs().write_escaped(F.getName()) << '\n';
 return false;</pre>

}; // end of struct Hello
} // end of anonymous namespace

static RegisterStandardPasses Y(
 PassManagerBuilder::EP_EarlyAsPossible,
 [](const PassManagerBuilder &Builder,
 legacy::PassManagerBase &PM) { PM.add(new Hello()); });

```
43
  namespace -
44
     struct Hello2 : public FunctionPass {
45
46
       static char ID; // Pass identification, replacement for typeid
47
       Hello2() : FunctionPass(ID) {}
48
49
50
       bool runOnFunction(Function &F) override {
           unsigned int basicBlockCount = 0;
                                                          Here is where we will accumulate the basic
51
           unsigned int instructionCount = 0;
                                                           blocks and instructions within our function
52
           for(BasicBlock & bb : F){
53
54
                ++basicBlockCount;
                for(Instruction &i : bb){
55
56
57
                    ++instructionCount;
58
59
           errs() << "Hello2 is running: ";</pre>
           errs().write escaped(F.getName())
                                                  << "Basic Blocks:" << basicBlockCount</pre>
60
61
62
63
                                                  << "Instruction:" << instructionCount << "\n"
64
       void getAnalysisUsage(AnalysisUsage &AU) const override {
65
         AU.setPreservesAll();
66
67
     };
68
69
70 char Hello2::ID = 0;
71 static RegisterPass<Hello2>
72 Y("hello2", "Hello World Pass (with getAnalysisUsage implemented)");
```

```
74 #include "llvm/IR/CallSite.h"
75 namespace {
76
     struct Hello3 : public FunctionPass {
77
78
       static char ID; // Pass identification, replacement for typeid
79
       Hello3() : FunctionPass(ID) {}
80
81
       bool runOnFunction(Function &F) override {
82
            for(BasicBlock &bb: F){
83
                for(Instruction &i: bb){
84
85
                    CallSite cs(&i);
86
                    if(!cs.getInstruction()){
87
                        continue;
88
89
                    Value *called = cs.getCalledValue()->stripPointerCasts();
90
                    if(Function* f = dyn cast<Function>(called)){
91
                        errs() << "\tDirect call to function:" << f->getName()
92
                                << " from "
                                                                 << F.getName() << "\n";
93
94
95
96
97
            return false;
98
99
100
101
       void getAnalysisUsage(AnalysisUsage &AU) const override {
102
         AU.setPreservesAll();
103
104
105 }
106
107 char Hello3::ID = 0;
LO8 static RegisterPass<Hello3>
L09 Z("hello3", "Hello World Pass (Get direct calls)");
```

The Module pass

```
bool runOnModule(Module &M) override {
    setupHooks(" Z10 initMaini",M);
    Module::FunctionListType &functions = M.getFunctionLi
    for(Module::FunctionListType::iterator FI = function
        if(" Z10=initMaini"==FI->getName()){
            continue;
        if("main"==FI->getName()){
            InstrumentEnterFunction(" Z10 initMaini",*FI,M);
        }
    return true;
```

LLVM Dataflow Analysis

• LLVM IR is in SSA form:

- use-def and def-use chains are always available
- All objects have user/use info, even functions
- Control Flow Graph is always available:
 - Exposed as BasicBlock predecessor/successor lists
 - Many generic graph algorithms usable with the CFG
- Higher-level info implemented as passes:
 - Dominators, CallGraph, induction vars, aliasing, GVN, ...

See also: docs/ProgrammersManual.html

Homework

- Read "llvm: a compilation framework for lifelong program analysis & transformation" 2014
- Read "SoK: Sanitizing for Security". Oakland 2019